

GLOBAL RIGIDITY OF GENERIC FRAMEWORKS ON THE CYLINDER

BILL JACKSON AND ANTHONY NIXON

ABSTRACT. We show that a generic framework (G, p) on the cylinder is globally rigid if and only if G is a complete graph on at most four vertices or G is both redundantly rigid and 2-connected. To prove the theorem we also derive a new recursive construction of circuits in the simple $(2, 2)$ -sparse matroid, and a characterisation of rigidity for generic frameworks on the cylinder when a single designated vertex is allowed to move off the cylinder.

1. INTRODUCTION

A (bar-joint) framework (G, p) in \mathbb{R}^d is the combination of a finite, simple graph $G = (V, E)$ and a realisation $p : V \rightarrow \mathbb{R}^d$. The framework (G, p) is rigid if every edge-length preserving continuous motion of the vertices arises as a congruence of \mathbb{R}^d . Moreover (G, p) is globally rigid if every framework (G, q) with the same edge lengths as (G, p) arises from a congruence of \mathbb{R}^d .

In general it is NP-hard to determine the rigidity or global rigidity of a given framework [1, 22]. These problems become more tractable, however, for *generic frameworks in \mathbb{R}^d* , i.e. frameworks whose coordinates are algebraically independent over \mathbb{Q} . It is known that both the rigidity and global rigidity of a generic framework (G, p) in \mathbb{R}^d depend only on the underlying graph G , see [2, 8]. We say that G is *rigid* or *globally rigid in \mathbb{R}^d* if some/every generic realisation of G in \mathbb{R}^d has the corresponding property. Combinatorial characterisations of generic rigidity and global rigidity in \mathbb{R}^d have been obtained when $d \leq 2$, see [15, 10], and these characterisations give rise to efficient combinatorial algorithms to decide if these properties hold. In higher dimensions, however, no combinatorial characterisations or algorithms are yet known.

We consider the situation where (G, p) is a framework in \mathbb{R}^3 whose vertices are constrained to lie on a fixed surface. Combinatorial characterisations for generic rigidity in this context were established for the sphere [24, 19], cylinder [19], and cone [20]. In particular it was shown that a generic realisation of a graph G on the sphere is rigid if and only if G is rigid in the plane.

The characterisation of rigidity for a generic framework (G, p) on the cylinder uses the *simple $(2, 2)$ -sparse matroid*. This is the matroid $\mathcal{M}_{2,2}^*$ on the edge set of a large complete graph K_n in which a set of edges F is *independent* if and only if $|F'| \leq 2|V(F')| - 2$ for all $\emptyset \neq F' \subseteq F$, with strict inequality when $|F'| = 2$. For $G \subseteq K_n$, we denote the submatroid of $\mathcal{M}_{2,2}^*$ induced on $E(G)$ by $\mathcal{M}_{2,2}^*(G)$.

Theorem 1.1. *Let (G, p) be a generic framework on the cylinder. Then (G, p) is rigid if and only if G is a complete graph on at most 3 vertices or $\mathcal{M}_{2,2}^*(G)$ has rank $2|V(G)| - 2$.*

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This characterisation implies that the rigidity of a generic framework (G, p) on a cylinder depends only on the underlying graph G . We will say that a graph G is *rigid on the cylinder* if it is generically rigid on the cylinder and that G is *redundantly rigid on the cylinder* if $G - e$ is rigid on the cylinder for all edges e of G .

We next consider global rigidity on surfaces. Such problems arise naturally as sensor network localisation problems, see [17, Page 2] and the references therein. Connelly and Whiteley [6] showed that a graph G is generically globally rigid on the sphere if and only if it is generically globally rigid in the plane (which holds if and only if G is 3-connected and redundantly rigid in the plane by [10]). In [13], necessary combinatorial conditions were established for a framework on a surface to be generically globally rigid. The conditions, redundant rigidity and k -connectivity (where the integer k depends on the chosen surface), are analogous to those which characterise generic global rigidity on the plane and the sphere. These conditions were conjectured to also be sufficient for generic frameworks on cylinders and cones. In this paper we verify this conjecture for the cylinder.

Theorem 1.2. *Let (G, p) be a generic framework on the cylinder. Then (G, p) is globally rigid on the cylinder if and only if G is either a complete graph on at most four vertices or G is 2-connected and redundantly rigid on the cylinder.*

The key step in proving sufficiency in Theorem 1.2 is the following result which deals with the special case when G is 2-connected and redundantly rigid with the minimum possible number of edges. Theorem 1.1 implies that $|E| \geq 2|V| - 1$ whenever $G = (V, E)$ is redundantly rigid, and that equality holds only if E is a circuit in $\mathcal{M}_{2,2}^*$. We will abuse terminology and say that G is a *circuit in $\mathcal{M}_{2,2}^*$* whenever this occurs. For simplicity we will refer to such graphs as $\mathcal{M}_{2,2}^*$ -circuits.

Theorem 1.3. *Let G be an $\mathcal{M}_{2,2}^*$ -circuit and (G, p) be a generic framework on the cylinder. Then (G, p) is globally rigid on the cylinder.*

We will need the following three results to prove Theorem 1.3. The first is a decomposition result for $\mathcal{M}_{2,2}^*$ -circuits which uses the graph operations defined in Figure 1.

Theorem 1.4. [18, Lemmas 3.1, 3.2, 3.3] *Suppose G_0, G_1, G_2 are graphs with $|E(G_i)| = 2|V(G_i)| - 1$ for all $0 \leq i \leq 2$ and that G_0 is a j -join of G_1 and G_2 for some $1 \leq j \leq 3$. Then G_0 is a $\mathcal{M}_{2,2}^*$ -circuit if and only if both G_1 and G_2 are $\mathcal{M}_{2,2}^*$ -circuits.*

The second result we shall need is a recursive construction for $\mathcal{M}_{2,2}^*$ -circuits which uses the j -join operations as well as the *1-extension operation* which deletes an edge xy from a graph G and then adds a new vertex v and three new edges vx, vy, vz for some vertex $z \neq x, y$. The recursion begins with the three $\mathcal{M}_{2,2}^*$ -circuits defined in Figure 2.

Theorem 1.5. [18, Theorem 1.1] *Suppose G is an $\mathcal{M}_{2,2}^*$ -circuit. Then G can be obtained from either K_5^- , H_1 or H_2 by recursively applying the operations of 1-extension, and 1-, 2- and 3-join.*

The third result tells us that two equivalent generic frameworks on the cylinder have closely related equilibrium stresses (which will be defined in Section 3).

Theorem 1.6. [14, Theorem 12] *Let (G, p) be a generic framework on the cylinder and (ω, λ) be an equilibrium stress for (G, p) . Let (G, q) be equivalent to (G, p) . Then (ω, λ') is an equilibrium stress for (G, q) for some $\lambda' \in \mathbb{R}^n$.*

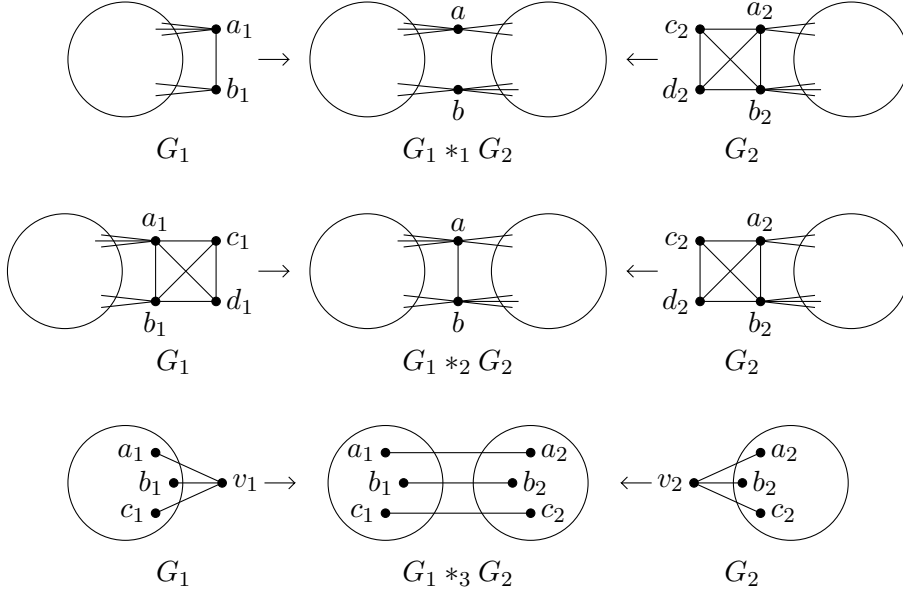


FIGURE 1. The 1-, 2- and 3-join operations. The 1- and 2-join operations form the graphs in the centre by merging a_1 and a_2 into a , and b_1 and b_2 into b . We write $G = G_1 *_j G_2$ to mean G is a j -join of G_1 and G_2 .

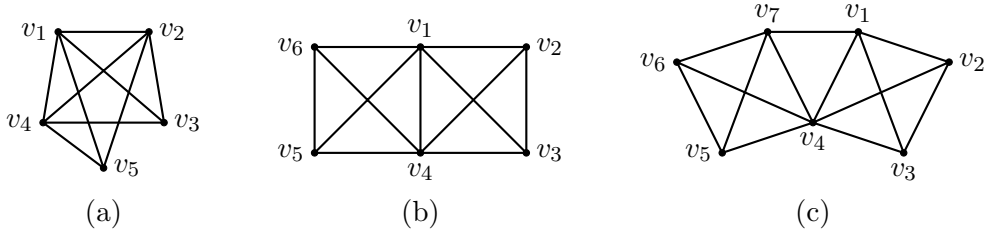


FIGURE 2. The graphs K_5^- , H_1 and H_2 .

An overview of the proof of Theorems 1.2 and 1.3 is as follows. A key step in showing that 2-connectivity and redundant rigidity are sufficient to imply generic global rigidity for $\mathcal{M}_{2,2}^*$ -circuits is to show inductively that every generic realisation of an $\mathcal{M}_{2,2}^*$ -circuit on the cylinder has a maximum rank equilibrium stress. The most straightforward way to do this would be to show that the operations used in Theorem 1.5 preserve the property of having such a stress, but we are unable to show that the join operations do this. Instead we use Theorem 1.5 to obtain a new recursive construction which uses operations which we can show preserve this property. We derive this construction in Section 2, give formal definitions for infinitesimal rigidity and equilibrium stresses in Section 3, and show that each of the operations in our recursive construction preserves the stress property in Section 4. A technical detail in the proof in Section 4, that the equilibrium stress is nowhere zero, is dealt with in Section 6 and gives rise to a characterisation of rigidity for generic frameworks on the cylinder in which one designated vertex is allowed to move freely in \mathbb{R}^3 .

Given the result that every generic realisation of an $\mathcal{M}_{2,2}^*$ -circuit has a maximum rank equilibrium stress, we can apply Theorem 1.6 to deduce that any two equivalent generic realisations of an $\mathcal{M}_{2,2}^*$ -circuit on the cylinder have related maximum rank equilibrium stresses.

If we could show that the stresses were identical then we could immediately deduce that the two realisations are congruent using [14, Theorem 9]. Since we cannot show this, we instead use the above mentioned relation between the two stresses to first deduce that the two frameworks are ‘VR-equivalent’ i.e. the projections of the vertices of the frameworks onto the axis of the cylinder are linked by a dilation. We then apply a characterization of ‘global VR-rigidity’ (which we subsequently derive in Section 7) to deduce that the two realisations are congruent. This gives Theorem 1.3. We then use an ‘ear-decomposition’ of the rigidity matroid to extend the characterization to arbitrary graphs and obtain Theorem 1.2. Both characterisations are derived in Section 5. We close with a short section which points out that the property of having a maximum rank stress matrix is a sufficient condition for generic global rigidity on the cylinder (a result we had previously conjectured in [14, Conjecture 1]), and that our characterization of global rigidity on the cylinder also holds for a generic framework in \mathbb{R}^3 with the added constraints that the distances of the vertices from a given line are fixed.

2. RECURSIVE CONSTRUCTION

We will refine the recursive construction for $\mathcal{M}_{2,2}^*$ -circuits given in Theorem 1.5. The aim is to make the steps in the recursion as simple as possible since we have to show they preserve global rigidity, or, more precisely, preserve the property of having a maximum rank equilibrium stress. To this end we introduce two ‘new’ operations. Given an $\mathcal{M}_{2,2}^*$ -circuit $G = (V, E)$, the first operation, K_4^- -extension, is just a 1-join of G and H_1 . The second operation, *generalised vertex split*, is defined as follows: choose $v \in V$ and a partition N_1, N_2 of the neighbours of v ; then delete v from G and add two new vertices v_1, v_2 joined to N_1, N_2 , respectively; finally add two new edges v_1v_2, v_1x for some $x \in V \setminus N_1$. These operations are illustrated in Figures 3 and 4.

The usual vertex splitting operation, see [25], is the special case when x is chosen to be a neighbour of v_2 . Note also that the special case when v_1 has degree 3 (and $v_2 = v$) is the 1-extension operation. At times it will be convenient to work directly with the 1-extension operation itself, for example in Theorem 2.2 below.

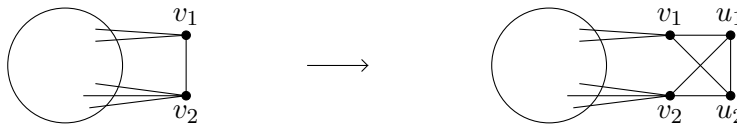


FIGURE 3. K_4^- -extension.

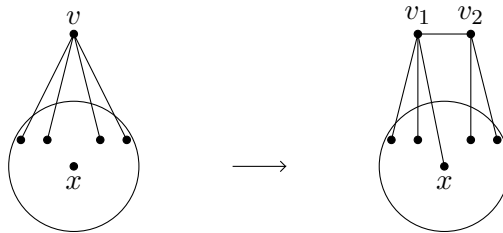


FIGURE 4. Generalised vertex split.

We can now state our simplified recursive construction for $\mathcal{M}_{2,2}^*$ -circuits.

Theorem 2.1. *Suppose G is an $\mathcal{M}_{2,2}^*$ -circuit. Then G can be obtained from either K_5^- or H_1 by recursively applying the operations of K_4^- -extension and generalised vertex split, in such a way that each of the intermediate graphs is an $\mathcal{M}_{2,2}^*$ -circuit.*

Since the K_4^- -extension operation is a special case of the 1-join operation, it must necessarily preserve the property of being an $\mathcal{M}_{2,2}^*$ -circuit by Theorem 1.4. The 1-extension operation also preserves the property of being an $\mathcal{M}_{2,2}^*$ -circuit by [18, Lemma 2.1], but the generalised vertex split operation may not.

We will refer to the inverse operations to those used in Theorem 2.1 as K_4^- -reduction and edge-reduction, respectively, and to the inverse of the 1-extension operation as 1-reduction. (Thus edge-reduction chooses two adjacent edges e, f and then deletes e and contracts f .) We say that an application of these operations is *admissible* if, when we apply it to an $\mathcal{M}_{2,2}^*$ -circuit, we obtain a smaller $\mathcal{M}_{2,2}^*$ -circuit. Theorem 1.4 implies that K_4^- -reduction will always be admissible but the operations of 1-reduction and, more generally, edge-reduction may not.

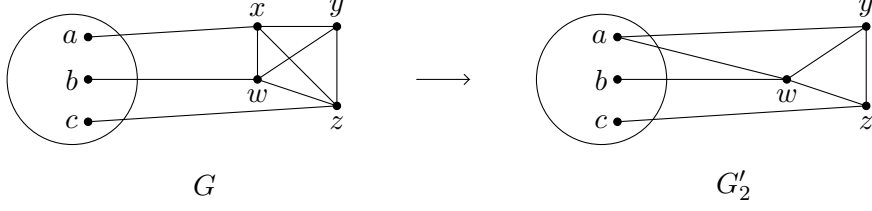
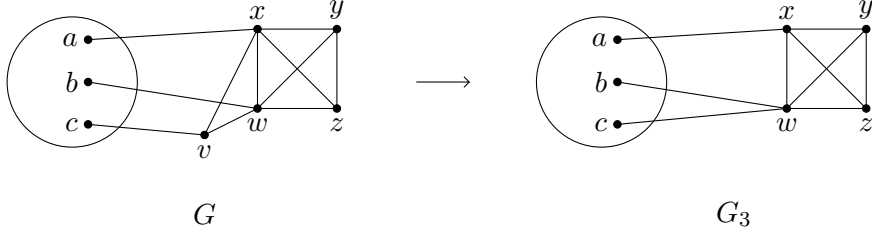
We say that a vertex is a *node* of a graph if it has degree three, and that a node v in an $\mathcal{M}_{2,2}^*$ -circuit G is *admissible* if we can construct a smaller $\mathcal{M}_{2,2}^*$ -circuit from G by applying a 1-reduction operation at v . A graph is *essentially 4-edge-connected* if it is 3-edge-connected and the only edge cuts of size 3 consist of 3 edges incident to a node. It was shown in [18, Theorem 1.2] that every 3-connected, essentially 4-edge-connected $\mathcal{M}_{2,2}^*$ -circuit other than K_5^- has at least two admissible nodes. We first extend this result by allowing certain 2-vertex-cuts and 3-edge-cuts. We need the following definitions.

A *2-vertex-separation* of a graph G is a pair of induced subgraphs $F_1 = (V_1, E_1)$ and $F_2 = (V_2, E_2)$ such that $F_1 \cup F_2 = G$, $|V_1 \cap V_2| = 2$ and $V_1 \setminus V_2 \neq \emptyset \neq V_2 \setminus V_1$. The 2-separation (F_1, F_2) is *nontrivial* if $F_i \neq K_4$ for each $i = 1, 2$. A *3-edge-separation* is a pair of vertex-disjoint induced subgraphs (F_1, F_2) such that $F_1 \cup F_2 = G - S$ for some set $S \subseteq E$ with $|S| = 3$. It is *nontrivial* if S is a set of three independent edges in G . An *atom* of G is a subgraph F such that F is an element of a nontrivial 2-vertex separation or 3-edge-separation of G and no proper subgraph of F has this property.

We also need to extend the concepts of circuits and admissible nodes to multigraphs. The $(2, 2)$ -sparse matroid $\mathcal{M}_{2,2}$ is the matroid on the edge set of a large complete multigraph in which a set of edges F is independent if $|F'| \leq 2|V(F')| - 2$ for all $\emptyset \neq F' \subseteq F$. Suppose that H is an $\mathcal{M}_{2,2}$ -circuit i.e. $E(H)$ is a circuit in $\mathcal{M}_{2,2}$. We say that a node v of H is *allowable* if applying the 1-reduction operation at v produces a smaller $\mathcal{M}_{2,2}$ -circuit and does not create any new multiple edges.

Theorem 2.2. *Suppose G is an $\mathcal{M}_{2,2}^*$ -circuit which is distinct from K_5^- , H_1 and H_2 , and that G has no nontrivial 2-vertex separation and no nontrivial 3-edge-separation. Then G has at least two admissible nodes.*

Proof. If G is 3-connected then the statement is [18, Theorem 1.2] so we may assume that G is not 3-connected. Since G is an $\mathcal{M}_{2,2}^*$ -circuit, G is 2-connected by [18, Lemma 2.3]. Since G has no nontrivial 2-vertex separation and $G \neq H_1$, every 2-vertex-separation (F_i, F_j) of G has $F_i = K_4$ and $F_j \neq K_4$. For every such 2-vertex-separation with $V(F_i) \cap V(F_j) = \{x_i, y_i\}$, we consider the multigraph H formed by deleting $V(F_i) - \{x_i, y_i\}$ and $E(F_i)$ from G , and adding two copies of the edge $x_i y_i$. Since $G \neq H_2$, H is a 3-connected circuit in $\mathcal{M}_{2,2}$ with no nontrivial 3-edge-separation. Also, each x_i and y_i has degree at least 4 in H ; otherwise

FIGURE 5. The subcase $G_1 = K_5^-$ in Case 1.FIGURE 6. The subcase $G_1 = H_1$ of Case 1.

there would be a nontrivial 2-vertex-separation in G . It follows that every allowable node in H is an admissible node of G . We can now show that H contains two allowable nodes using the proof technique of [18, Theorem 1.2], see [18, Section 4]. \square

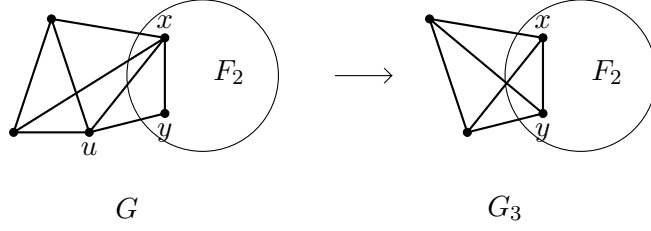
Theorem 2.1 will follow immediately from the following reduction result.

Theorem 2.3. *Let $G = (V, E)$ be an $\mathcal{M}_{2,2}^*$ -circuit distinct from K_5^- and H_1 . Then G has either a K_4^- -reduction or an admissible edge-reduction.*

Proof. We proceed by induction on $|V|$. If G has no nontrivial 2-separation or 3-edge-separation then either $G = H_2$ and can be reduced to H_1 by an admissible edge-reduction operation, or else G has two admissible nodes by Theorem 2.2, and hence has two admissible edge-reductions. Thus we may suppose that this is not the case. It follows that G has at least two distinct atoms. Let F_1 be an atom of G and (F_1, F_2) be the nontrivial separation which contains F_1 . Consider the following cases.

Case 1: (F_1, F_2) is a nontrivial 3-edge-separation. Let G_i be obtained from G by contracting F_{3-i} to a single vertex z_i for each $i \in \{1, 2\}$. Then $G = G_1 *_3 G_2$ so G_1, G_2 are $\mathcal{M}_{2,2}^*$ -circuits by Theorem 1.4. Since F_1 is an atom, G_1 has no nontrivial 2-vertex-separation or 3-edge-separation. If $G_1 \notin \{K_5^-, H_1, H_2\}$ then G_1 has an admissible node v distinct from z_1 by Theorem 2.2, and v will be an admissible node in G by Theorem 1.4.

If $G_1 = K_5^-$ then G is as shown in Figure 5. We can apply an edge-reduction operation which deletes the edge xw and then contracts the edge xa to construct the graph G'_2 on the right of Figure 5. Then G'_2 is an $\mathcal{M}_{2,2}^*$ -circuit since it can be obtained from G_2 by two 1-extensions. Hence this edge-reduction operation is admissible. If $G_1 = H_2$ then we will contradict the assumption that F_1 is an atom of G . It remains to consider the subcase when $G_1 = H_1$, which is illustrated in Figure 6. In this case the vertex v is admissible, since performing a 1-reduction on v gives the graph G_3 on the right of Figure 6, and we have $G_3 = G_2 *_1 H_2$ so G_3 is an $\mathcal{M}_{2,2}^*$ -circuit by Theorem 1.4.

FIGURE 7. The graphs G and G_3 in the subcase $G_1 = H_2$ of Case 2.

Case 2: (F_1, F_2) is a nontrivial 2-vertex-separation, $V(F_1) \cap V(F_2) = \{x, y\}$ and $E(F_1) \cap E(F_2) = \{xy\}$. Let G_i be obtained from F_i by adding two new vertices $\{w, z\}$ and five new edges $\{wx, wy, wz, xz, yz\}$. Then $G = G_1 *_2 G_2$ so G_1 and G_2 are $\mathcal{M}_{2,2}^*$ -circuits by Theorem 1.4. Since F_1 is an atom, G_1 has no nontrivial 2-vertex-separation or 3-edge-separation. We have $G_1 \neq K_5^-$ since G_1 is not 3-connected and $G_1 \neq H_1$ since the 2-separation (F_1, F_2) is nontrivial.

Suppose $G_1 = H_2$. Then G is as shown in Figure 7. Let G_3 be obtained from G by the edge-reduction which deletes xy and then contracts uy . This an admissible edge-reduction since $G = H_2 *_2 G_3$ so G_3 is an $\mathcal{M}_{2,2}^*$ -circuit by Theorem 1.4.

Hence we may suppose that $G_1 \notin \{K_5^-, H_1, H_2\}$. Then G_1 has an admissible node v by Theorem 2.2. The vertex v will be distinct from w, z since they are not admissible, and distinct from x, y since they are not nodes in G_1 . Thus v will be an admissible node in G .

Case 3: (F_1, F_2) is a nontrivial 2-vertex-separation, $V(F_1) \cap V(F_2) = \{x, y\}$ and $xy \notin E$. We have

$$|E(F_1)| + |E(F_2)| = |E(G)| = 2|V(G)| - 1 = 2|V(F_1)| + 2|V(F_2)| - 5$$

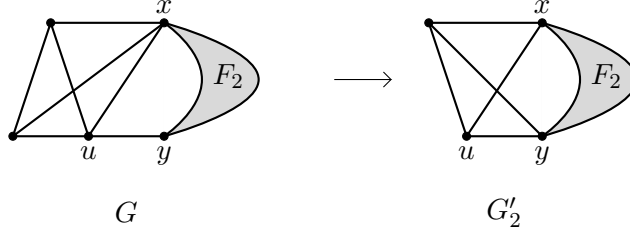
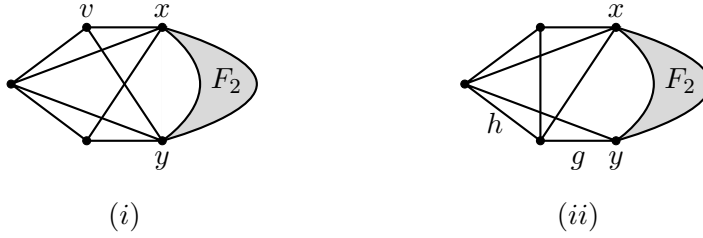
and $|E(F_i)| \leq 2|V(F_i)| - 2$ for each $i = 1, 2$ so $2|V(F_i)| - 3 \leq |E(F_i)| \leq 2|V(F_i)| - 2$. Consider the following subcases.

Subcase 3.1: $|E(F_1)| = 2|V(F_1)| - 3$. Let G_1 be obtained from F_1 by adding two new vertices $\{w, z\}$ and six new edges $\{wx, wy, wz, xy, xz, yz\}$, and $G_2 = F_2 + xy$. Then $G = G_1 *_1 G_2$ so G_1 and G_2 are $\mathcal{M}_{2,2}^*$ -circuits by Theorem 1.4. Since F_1 is an atom, G_1 has no nontrivial 2-vertex-separation or 3-edge-separation. We have $G_1 \neq K_5^-$ since G_1 is not 3-connected. If $G_1 = H_1$ then G_2 is a K_4^- -reduction of G . Hence we may assume that $G_1 \neq H_1$.

Suppose $G_1 = H_2$. Then G is as shown in Figure 8. Let G'_2 be obtained from G by the edge-reduction which deletes xu and then contracts uy . This an admissible edge-reduction since $G'_2 = H_1 *_1 G_2$ so G'_2 is an $\mathcal{M}_{2,2}^*$ -circuit by Theorem 1.4.

Hence we may suppose that $G_1 \notin \{K_5^-, H_1, H_2\}$. Then G_1 has an admissible node v by Theorem 2.2. The vertex v will be distinct from w, z since they are not admissible, and distinct from x, y since they are not nodes in G_1 . Thus v will be an admissible node in G .

Subcase 3.2: $|E(F_1)| = 2|V(F_1)| - 2$. Let G_2 be obtained from F_2 by adding two new vertices $\{w, z\}$ and six new edges $\{wx, wy, wz, xy, xz, yz\}$, and $G_1 = F_1 + xy$. Then $G = G_1 *_1 G_2$ so G_1 and G_2 are $\mathcal{M}_{2,2}^*$ -circuits by Theorem 1.4. We may apply the argument used in Subcase 3.1 (to show that $G_1 \notin \{K_5^-, H_1, H_2\}$) to deduce that $G_2 \notin \{K_5^-, H_1, H_2\}$.

FIGURE 8. The graphs G and G'_2 in the subcase $G_1 = H_2$ of Subcase 3.1.FIGURE 9. The graph G in the subcase $G_1 = K_5^-$ of Subcase 3.2.

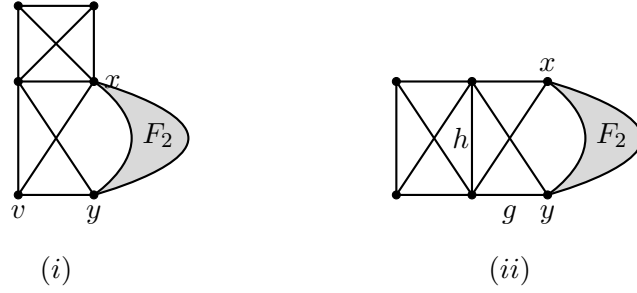
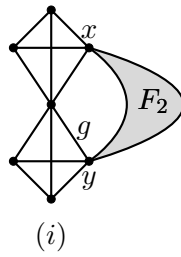
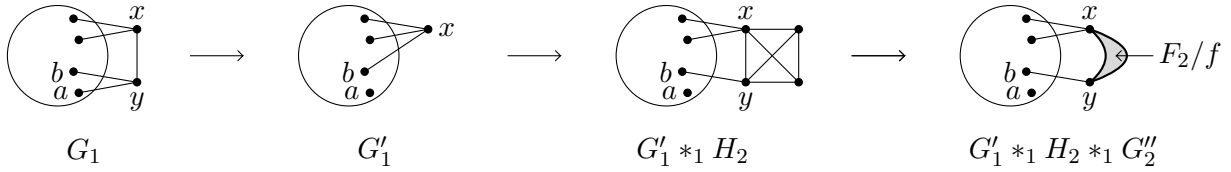
By induction, G_2 has either a K_4^- -reduction or an admissible edge-reduction. If G_2 has a K_4^- -reduction, then the same reduction will exist in G . Hence we may suppose that G_2 has an admissible edge-reduction which deletes an edge e then contracts an adjacent edge f . If $e, f \in E(F_2)$, then the same edge-reduction will be admissible in G , so we may assume that this is not the case. Since the edge-reduction is admissible in G_2 we must have $f \notin \{wx, wy, wz, xy, xz, yz\}$, $e = xy$, and f is incident to either x or y . Without loss of generality we may suppose that f is incident with y . Let G'_2 be the $\mathcal{M}_{2,2}^*$ -circuit created by this edge-reduction and let G''_2 be the graph obtained from G'_2 by deleting the vertices w, z and adding the edge xy . Since $G'_2 = G''_2 *_1 H_1$, G''_2 is an $\mathcal{M}_{2,2}^*$ -circuit by Theorem 1.4.

We next consider G_1 . Since F_1 is an atom, G_1 has no nontrivial 2-vertex-separation or 3-edge-separation.

Suppose $G_1 = K_5^-$. Then G has the structure of one of the graphs shown in Figure 9. If case (i) occurs then the 1-reduction of G which deletes v and adds the edge xy gives the $\mathcal{M}_{2,2}^*$ -circuit G_2 , so is admissible. If case (ii) occurs then the edge-reduction of G which deletes h and contracts g also gives the $\mathcal{M}_{2,2}^*$ -circuit G_2 , so is admissible.

Suppose $G_1 = H_1$. Since F_1 is an atom, G would have the structure of one of the graphs shown in Figure 10. If case (i) occurs then the 1-reduction of G which deletes v and adds the edge xy gives the $\mathcal{M}_{2,2}^*$ -circuit $G_2 *_2 H_2$, so is admissible. If case (ii) occurs then the edge-reduction of G which deletes h and contracts g also gives the $\mathcal{M}_{2,2}^*$ -circuit $G_2 *_2 H_2$, so is admissible.

Suppose $G_1 = H_2$. Since F_1 is an atom, G would have the structure of the graph shown in Figure 11(i). We can now use the admissible edge-reduction of G_2 to G'_2 described in the first paragraph of this subcase. Consider the edge-reduction of G which deletes the edge g shown in Figure 11(i) and contracts the edge $f \in E(F_2)$. This is admissible since it gives the $\mathcal{M}_{2,2}^*$ -circuit $G''_2 *_1 H_3$, where H_3 is the $\mathcal{M}_{2,2}^*$ -circuit shown in Figure 11(ii). (Note that H_3 is an $\mathcal{M}_{2,2}^*$ -circuit since $H_3 = H_1 *_1 H_2$.)


 FIGURE 10. The graph G in the subcase $G_1 = H_1$ of Subcase 3.2.

 FIGURE 11. The graphs G and H_3 in the subcase $G_1 = H_2$ of Subcase 3.2.

 FIGURE 12. The subcase of Subcase 3.2 when $G_1 \notin \{K_5^-, H_1, H_2\}$.

Hence we may assume that $G_1 \notin \{K_5^-, H_1, H_2\}$. Then G_1 has two admissible nodes by Theorem 2.2. It can be seen that any admissible node of G_1 which is distinct from x, y is an admissible node of G , so we may suppose that x and y are the only admissible nodes of G_1 . Then both x and y have degree three in G_1 . Let a, b, x be the neighbours of y in G_1 . Since x has degree two in $G_1 - y$, the admissible 1-reduction at y constructs a new graph G'_1 from G_1 by deleting y and then adding a new edge from x to either a or b , say b . We use the admissible reduction of G_2 to G''_2 described in the first paragraph of this subcase. Let G' be the graph obtained from G by performing an edge-reduction which deletes the edge ay and then contracts the edge $f \in E(F_2)$. Then $G' = G'_1 *_1 H_2 *_1 G''_2$ and hence G' is an $\mathcal{M}_{2,2}^*$ -circuit, see Figure 12. Thus the edge-reduction which transforms G to G' is admissible. \square

To see that the K_4^- -extension operation is needed in Theorem 2.3, observe that we can construct graphs which do not admit admissible edge-reductions as follows. Take any $\mathcal{M}_{2,2}^*$ -circuit H , and apply the K_4^- -extension to every single edge of G . The resulting graph G has two types of edges. Those edges with no end-vertices in H are contained in two triangles so any edge-reduction which contracts such an edge results in a non-simple graph. The remaining edges, those with exactly one end-vertex in H , are not admissible either since

any edge-reduction which contracts such an edge results in a graph containing a vertex of degree two.

3. RIGIDITY AND STRESS MATRICES

Let $G = (V, E)$ be a graph with $V = \{v_1, \dots, v_n\}$. We will consider realisations of G on the unit cylinder $\mathcal{Y} = \{(x, y, z) \in \mathbb{R}^3 : x^2 + y^2 = 1\}$.¹ A *framework* (G, p) on \mathcal{Y} is an ordered pair consisting of a graph G and a realisation p such that $p(v_i) \in \mathcal{Y}$ for all $v_i \in V$.

Two frameworks (G, p) and (G, q) on \mathcal{Y} are *equivalent* if $\|p(v_i) - p(v_j)\| = \|q(v_i) - q(v_j)\|$ for all edges $v_i v_j \in E$. Moreover (G, p) and (G, q) are *congruent* if $\|p(v_i) - p(v_j)\| = \|q(v_i) - q(v_j)\|$ for all pairs of vertices $v_i, v_j \in V$. The framework (G, p) is *globally rigid* on \mathcal{Y} if every equivalent framework (G, q) on \mathcal{Y} is congruent to (G, p) . It is *rigid* on \mathcal{Y} if there exists an $\epsilon > 0$ such that every framework (G, q) on \mathcal{Y} which is equivalent to (G, p) , and has $\|p(v_i) - q(v_i)\| < \epsilon$ for all $1 \leq i \leq n$, is congruent to (G, p) . It is *generic* on \mathcal{Y} if $\text{td}[\mathbb{Q}(p) : \mathbb{Q}] = 2n$.

The *rigidity matrix* $R_{\text{cyl}}(G, p)$ of a framework (G, p) on \mathcal{Y} is the $(|E| + |V|) \times 3|V|$ matrix

$$R_{\text{cyl}}(G, p) = \begin{pmatrix} R(G, p) \\ S(G, p) \end{pmatrix}$$

where: $R(G, p)$ has rows indexed by E and 3-tuples of columns indexed by V in which, for $e = v_i v_j \in E$, the submatrices in row e and columns v_i and v_j are $p(v_i) - p(v_j)$ and $p(v_j) - p(v_i)$, respectively, and all other entries are zero; $S(G, p)$ has rows indexed by V and 3-tuples of columns indexed by V in which, for $v_i \in V$, the submatrix in row v_i and column v_i is $\bar{p}(v_i) = (x_i, y_i, 0)$ when $p(v_i) = (x_i, y_i, z_i)$. We refer to the vectors in the kernel of $R_{\text{cyl}}(G, p)$ as *infinitesimal flexes* of (G, p) . We will consider an infinitesimal flex as a map $s : V \rightarrow \mathbb{R}^3$ such that $s(v_i)$ is tangential to \mathcal{Y} at $p(v_i)$ for all $v_i \in V$ and $(p(v_j) - p(v_i)) \cdot (s(v_j) - s(v_i)) = 0$ for all $v_i v_j \in E$. The kernel of $R_{\text{cyl}}(G, p)$ will always contain two linearly independent flexes corresponding to a translation along the axis of the cylinder and a rotation about the same axis. This implies that $\text{rank } R_{\text{cyl}}(G, p) \leq 3n - 2$. We say that (G, p) is *infinitesimally rigid* on \mathcal{Y} if $\text{rank } R_{\text{cyl}}(G, p) = 3n - 2$. It was shown in [19] that a generic framework (G, p) on \mathcal{Y} is rigid if and only if it is a complete graph on at most three vertices or is infinitesimally rigid.

An *equilibrium stress* for a framework (G, p) on \mathcal{Y} is a pair (ω, λ) , where $\omega : E \rightarrow \mathbb{R}$ and $\lambda : V \rightarrow \mathbb{R}$ and (ω, λ) belongs to the cokernel of $R_{\text{cyl}}(G, p)$. Thus (ω, λ) is an equilibrium stress for (G, p) on \mathcal{Y} if and only if

$$(3.1) \quad \sum_{j=1}^n \omega_{ij}(p(v_i) - p(v_j)) + \lambda_i \bar{p}(v_i) = 0 \text{ for all } 1 \leq i \leq n,$$

where ω_{ij} is taken to be equal to ω_e if $e = v_i v_j \in E$ and to be equal to 0 if $v_i v_j \notin E$.

Given a stress (ω, λ) for a framework (G, p) on \mathcal{Y} we define $\Omega = \Omega(\omega)$ to be the $n \times n$ symmetric matrix with off-diagonal entries $-\omega_{ij}$ and diagonal entries $\sum_j \omega_{ij}$, and $\Lambda = \Lambda(\lambda)$ to be the $n \times n$ diagonal matrix with diagonal entries $\lambda_1, \lambda_2, \dots, \lambda_n$. The *stress matrix* associated to (ω, λ) is the $3n \times 3n$ symmetric matrix

$$\Omega_{\text{cyl}}(\omega, \lambda) = \begin{bmatrix} \Omega + \Lambda & 0 & 0 \\ 0 & \Omega + \Lambda & 0 \\ 0 & 0 & \Omega \end{bmatrix}.$$

¹For the purposes of (global) rigidity there is no loss in generality in assuming our cylinder has unit radius and is centred on the z -axis.

It follows from (3.1) that $(x_1, x_2, \dots, x_n), (y_1, y_2, \dots, y_n)$ belong to the cokernel of $\Omega + \Lambda$ and $(z_1, z_2, \dots, z_n), (1, 1, \dots, 1)$ belong to the cokernel of Ω . Hence $\text{rank}(\Omega + \Lambda) \leq n - 2$, $\text{rank} \Omega \leq n - 2$ and $\text{rank} \Omega_{\text{cyl}}(\omega, \lambda) \leq 3n - 6$. We will say that (ω, λ) has maximum rank when $\text{rank} \Omega_{\text{cyl}}(\omega, \lambda) = 3n - 6$.

4. MAXIMUM RANK EQUILIBRIUM STRESSES

We will show that every generic realisation of an $\mathcal{M}_{2,2}^*$ -circuit on the cylinder has a maximum rank equilibrium stress. We do this by showing that both of the recursive operations used in Theorem 2.1 preserve this property.

We will need the following elementary tool from linear algebra. Suppose $M = \begin{bmatrix} A & B \\ C & D \end{bmatrix}$ is a block matrix and A is invertible. Then the *Schur complement* of A in M is the matrix $F = D - CA^{-1}B$ and we have

$$(4.1) \quad \text{rank } M = \text{rank } A + \text{rank } F.$$

We will also need the following extension of [14, Lemmas 9 and 10].

Lemma 4.1. *Suppose (G, p) is an infinitesimally rigid framework on the cylinder \mathcal{Y} . Then (G, q) is infinitesimally rigid on \mathcal{Y} for all generic (G, q) on \mathcal{Y} . Moreover, if (ω, λ) is an equilibrium stress for (G, p) on \mathcal{Y} with $\text{rank} \Omega_{\text{cyl}}(\omega, \lambda) = 3n - 6$, then (G, q) has an equilibrium stress (ω', λ') on \mathcal{Y} with $\text{rank} \Omega_{\text{cyl}}(\omega', \lambda') = 3n - 6$ for all generic (G, q) on \mathcal{Y} .*

Proof. The first assertion follows from the fact that the rank $R_{\text{cyl}}(G, p)$ will be maximised at any generic framework (G, p) . We adapt the proof technique of Connelly and Whiteley [6, Theorem 5] to prove the second assertion. Choose an arbitrary infinitesimally rigid framework (G, p') on \mathcal{Y} . Since the entries in $R_{\text{cyl}}(G, p')$ are polynomials in p' and the space of equilibrium stresses of (G, p') is the cokernel of $R_{\text{cyl}}(G, p')$, each equilibrium stress of (G, p') can be expressed as a pair of rational functions $(\omega(p', t), \lambda(p', t))$ of p' and t , where t is a vector of $m - 2n + 2$ indeterminates. (To see this we can imagine solving $R_{\text{cyl}}(G, p)^T s = 0$ by Gaussian elimination. The resulting matrix will have $\text{rank} R_{\text{cyl}}(G, p) = 3n - 2$ leading entries. We can express a general solution by considering the stresses corresponding to the $(m+n) - (3n-2) = m - 2n + 2$ columns which do not contain leading entries as indeterminates and then solving for the stresses corresponding to the remaining columns.) This implies that the entries in the corresponding stress matrix $\Omega_{\text{cyl}}(\omega(p', t), \lambda(p', t))$ will also be rational functions of p' and t . Hence the rank of $\Omega_{\text{cyl}}(\omega(p', t), \lambda(p', t))$ will be maximised whenever (G, p') is generic on \mathcal{Y} and t is algebraically independent over $\mathbb{Q}(p')$. Hence, for any generic (G, q) on \mathcal{Y} , (G, q) is infinitesimally rigid on \mathcal{Y} and we can choose $t \in \mathbb{R}^{m-2n+2}$ such that $\text{rank} \Omega_{\mathcal{Y}}(\omega(q, t), \lambda(q, t)) = 3n - 6$. \square

We first show that the K_4^- -extension operation preserves the property of having a maximum rank equilibrium stress.

Theorem 4.2. *Suppose $G = (V, E)$ and $G_1 = (V_1, E_1)$ are graphs, $v_2v_3 \in E_1$, $G_1 - v_2v_3$ is rigid and G is a K_4^- -extension of G_1 on the edge v_2v_3 . Let (G, p) be a generic realisation of G on the unit cylinder \mathcal{Y} and let p_1 be the restriction of p to G_1 . Suppose (G_1, p_1) has a maximum rank equilibrium stress on \mathcal{Y} . Then (G, p) has a maximum rank equilibrium stress on \mathcal{Y} .*

Proof. Let $V_1 = \{v_2, v_3, \dots, v_n\}$ and suppose that G is constructed from G_1 by deleting v_2v_3 , adding two new vertices v_0, v_1 and five new edges $v_0v_1, v_0v_2, v_0v_3, v_1v_2, v_1v_3$. We may use the isometries of \mathcal{Y} to move (G, p_1) so that $p_1(v_2) = (0, 1, 0)$. Let $p_1(v_i) = (x_i, y_i, z_i)$

for $3 \leq i \leq n$. Define $q : V \rightarrow \mathbb{R}^3$ by putting $q(v) = p_1(v)$ for all $v \in V_1$ and choosing $q(v_0) = (0, 1, 1)$ and $q(v_1) = (-x_3, y_3, z_3)$. (We choose these values for q so that the rows of $R_{\text{cyl}}(G + v_2v_3, q)$ labelled by the vertices and edges of the subgraph H of $G + v_2v_3$ induced by $\{v_0, v_1, v_2, v_3\}$, are dependent. This will enable us to construct an equilibrium stress for (G, q) by combining equilibrium stresses for (G_1, p_1) and $(H, q|_H)$ in such a way that the net stress on v_2v_3 is zero.)

We first show that (G, q) is infinitesimally rigid on \mathcal{Y} . This follows from the facts that $G_1 - v_2v_3$ is rigid and that $(G - v_0v_1, q)$ can be constructed from $(G_1 - v_2v_3, p_1)$ by adding v_0 and v_1 as vertices of degree two at points which do not lie on the lines joining their two neighbours v_2, v_3 .

Now suppose (ω', λ') is a maximum rank equilibrium stress for (G_1, p_1) . Since $G_1 - v_2v_3$ is rigid, we may suppose (ω', λ') has been chosen so that the stress value on v_2v_3 is non-zero and hence we may scale (ω', λ') so that $\omega'_{23} = -1$. We may combine (ω', λ') with the unique equilibrium stress for $(H, q|_H)$ which has stress value one on v_2v_3 to obtain the equilibrium stress (ω, λ) for (G, q) on the unit cylinder defined by $\omega_f = \omega'_f$ for all $f \in E_1 - v_2v_3$, $\lambda_i = \lambda'_i$ for all $4 \leq i \leq n$, $\omega_{23} = 0$, $\omega_{12} = 1$, $\omega_{02} = -2z_3$, $\omega_{13} = \frac{1}{2y_3(z_3-1)}$, $\omega_{03} = -\frac{z_3}{z_3-1} = \omega_{01}$, $\lambda_2 = \lambda'_2 + 2y_3 - 2$, $\lambda_3 = \lambda'_3 + \frac{y_3-1}{y_3(z_3-1)}$, $\lambda_1 = \frac{y_3-1}{y_3(z_3-1)}$, and $\lambda_0 = -\frac{2z_3(y_3-1)}{z_3-1}$.

We have $\Omega(\omega) = \begin{bmatrix} A & B \\ B^T & D \end{bmatrix}$ where

$$A = \begin{bmatrix} \omega_{03} + \omega_{02} + \omega_{01} & -\omega_{01} \\ -\omega_{01} & \omega_{13} + \omega_{12} + \omega_{01} \end{bmatrix}, \quad B = \begin{bmatrix} -\omega_{02} & -\omega_{03} & 0 & \dots \\ -\omega_{12} & -\omega_{13} & 0 & \dots \end{bmatrix},$$

and

$$D = \begin{bmatrix} \sum_{j \geq 4} \omega_{2j} + \omega_{02} + \omega_{12} & 0 & -\omega_{24} & \dots \\ 0 & \sum_{j \geq 4} \omega_{3j} + \omega_{03} + \omega_{13} & -\omega_{34} & \dots \\ \vdots & \vdots & \vdots & \ddots \end{bmatrix}.$$

We can now substitute the values for $\omega_{01}, \omega_{02}, \omega_{03}, \omega_{12}, \omega_{13}$ into A, B to obtain

$$A = \begin{bmatrix} -2z_3 - \frac{2z_3}{z_3-1} & \frac{z_3}{z_3-1} \\ \frac{z_3}{z_3-1} & 1 + \frac{1}{2y_3(z_3-1)} - \frac{z_3}{z_3-1} \end{bmatrix} \quad \text{and} \quad B = \begin{bmatrix} 2z_3 & \frac{z_3}{z_3-1} & 0 & \dots \\ -1 & \frac{-1}{2y_3(z_3-1)} & 0 & \dots \end{bmatrix}.$$

Since $\{y_3, z_3\}$ is algebraically independent over \mathbb{Q} , A is invertible and a matrix calculation gives

$$B^T A^{-1} B = \begin{bmatrix} -2z_3 + 2 & -1 & 0 & \dots \\ -1 & -\frac{2y_3-1}{2y_3(z_3-1)} & 0 & \dots \\ 0 & 0 & 0 & \dots \\ \vdots & \vdots & \vdots & \ddots \end{bmatrix}.$$

We may now deduce that

$$D - B^T A^{-1} B = \begin{bmatrix} \sum_{j \geq 4} \omega_{2j} - 1 & 1 & -\omega_{24} & \dots \\ 1 & \sum_{j \geq 4} \omega_{3j} - 1 & -\omega_{34} & \dots \\ \vdots & \vdots & \vdots & \ddots \end{bmatrix} = \Omega(\omega').$$

Equation (4.1) now gives $\text{rank } \Omega(\omega) = \text{rank } A + \text{rank } \Omega(\omega') = |V(G)| - 2$.

We may use a similar calculation to deduce that $\text{rank}[\Omega(\omega) + \Lambda(\lambda)] = |V(G)| - 2$. Hence $\text{rank } \Omega_{\text{cyl}}(\omega, \lambda) = 3|V(G)| - 6$. The result now follows from Lemma 4.1. \square

We chose the values of $q(v_0)$ and $q(v_1)$ in the above proof so that the framework $(H, q|_H)$ would have a nowhere zero equilibrium stress. We could also have accomplished this by putting $q(v_0)$ and $q(v_1)$ in the same plane as $q(v_2)$ and $q(v_3)$, but such a choice would have resulted in the matrix A being singular.

We next consider the generalised vertex splitting operation. In [12, Theorem 5.2] it was proved that the standard vertex splitting operation, with the additional assumption that the new graph is rigid when we delete the *bridging edge* (i.e. the edge joining the two copies of the split vertex), preserves generic global rigidity on the cylinder. We will need a similar result for generalised vertex splitting and maximum rank equilibrium stresses. Our proof technique is based on a proof by Connelly [4] that the standard vertex splitting operation preserves the property of having a maximum rank equilibrium stress in Euclidean space. We first give a variant of [4, Lemma 28].

Lemma 4.3. *Let W_t be an $n \times n$ matrix whose entries are continuous functions of t for all $t \in (0, \delta)$ and X be a non-singular $n \times n$ matrix. Suppose that $X^T W_t X = \begin{bmatrix} A_t & B_t \\ C_t & D_t \end{bmatrix}$, that B_t, C_t, D_t tend to finite limits B, C, D , respectively, as t approaches 0, and that $A_t = (a_t)$ is a 1×1 matrix with $\lim_{t \rightarrow 0} |a_t| = \infty$. Then $\text{rank } W_t \geq \text{rank } D + 1$ for all t sufficiently close to 0.*

Proof. Equation (4.1) and the hypothesis that X is non-singular give

$$\text{rank } W_t = \text{rank } X^T W_t X = \text{rank } A_t + \text{rank}(D_t - a_t^{-1} C_t B_t).$$

Since $\lim_{t \rightarrow 0} |a_t| = \infty$ and $\lim_{t \rightarrow 0} D_t = D$, we have $\text{rank } A_t = 1$ and $\text{rank}(D_t - a_t^{-1} C_t B_t) \geq \text{rank } D$ when t is sufficiently close to zero. \square

We will also need the following result about frameworks with two coincident points. Let G be a graph with two distinguished vertices u and v . A framework (G, p) is *uv-coincident* if $p(u) = p(v)$. A *uv-coincident* framework is *uv-generic* if $(G - u, p')$ is generic, where p' is the restriction of p to $G - u$.

Theorem 4.4. [12, Theorem 18] *Let u and v be distinct vertices of a graph G and let (G, p) be a *uv-generic, uv-coincident* realisation of G on the cylinder. Then (G, p) is infinitesimally rigid if and only if the graphs $G - uv$ and G/uv are both rigid on the cylinder.*

We need one more result to show that the generalised vertex split operation preserves the property of having a maximum rank equilibrium stress when applied to $\mathcal{M}_{2,2}^*$ -circuits on \mathcal{Y} .

Lemma 4.5. *Let (G, p) be a generic realisation of an $\mathcal{M}_{2,2}^*$ -circuit on \mathcal{Y} and (ω, λ) be a non-zero equilibrium stress for (G, p) . Then $\lambda_i \neq 0$ for all $v_i \in V$.*

We will delay the proof of this lemma until Section 6 as its proof is fairly long and will be a distraction from our next result.

Theorem 4.6. *Suppose that (G, p) is a generic realisation of an $\mathcal{M}_{2,2}^*$ -circuit on \mathcal{Y} and that (ω, λ) is a maximum rank equilibrium stress for (G, p) . Let \hat{G} be obtained from G by a generalised vertex splitting operation and suppose that \hat{G} is an $\mathcal{M}_{2,2}^*$ -circuit on the cylinder. Then there exists a realisation (\hat{G}, q) on \mathcal{Y} which is infinitesimally rigid and has a maximum rank equilibrium stress.*

Proof. Suppose that $V(G) = \{v_1, v_2, \dots, v_n\}$ and that \hat{G} is obtained from G by choosing the vertex v_1 with neighbours v_2, v_3, \dots, v_m , deleting the edges $v_1 v_2, v_1 v_3, \dots, v_1 v_k$ and then adding a new vertex v_0 and new edges $v_0 v_1, v_0 v_2, \dots, v_0 v_k$ and $v_0 v_\ell$ for some $\ell \geq k + 1$.

Let (\hat{G}, \hat{p}) be the v_0v_1 -coincident v_0v_1 -generic framework with $\hat{p}|_G = p$ and $\hat{p}(v_0) = p(v_1)$. Then $(\hat{G} - v_0v_1, \hat{p})$ is infinitesimally rigid by Theorem 4.4 (since $\hat{G} - v_0v_1$ and $(\hat{G} - v_0v_1)/v_0v_1 = G + v_0v_\ell$ are both rigid on the cylinder). Hence $\text{rank } R_{\text{cyl}}(\hat{G} - v_0v_1, \hat{p}) = 3|V(\hat{G})| - 2$. Since \hat{G} is an $\mathcal{M}_{2,2}^*$ -circuit, this implies that the rows of $\hat{R}_{\text{cyl}}(\hat{G} - v_0v_1, \hat{p})$ are linearly independent.

Let $a = \sum_{j=2}^k \omega_{1j}(p(v_1) - p(v_j))$ and $b = \sum_{j=k+1}^m \omega_{1j}(p(v_1) - p(v_j))$. Since (ω, λ) is an equilibrium stress, we have $a + b + \lambda_1 \bar{p}(v_1) = 0$ and since $\lambda_1 \neq 0$ by Lemma 4.5, $\bar{p}(v_1) \in \text{span}\{a, b\}$.

We next show that $\dim \text{span}\{a, b\} = 2$. Suppose to the contrary that $\dim \text{span}\{a, b\} = 1$. Then a and b are both scalar multiples of $\bar{p}(v_1)$ so $a + \mu \bar{p}(v_1) = 0 = b + \nu \bar{p}(v_1)$ for some $\mu, \nu \in \mathbb{R}$. We can now define a non-zero equilibrium stress $(\hat{\omega}, \hat{\lambda})$ for $(\hat{G} - v_0v_1, \hat{p})$ as follows: $\hat{\omega}_{0j} = \omega_{1j}$ for $2 \leq j \leq k$ and otherwise $\hat{\omega}_{0j} = 0$, $\hat{\omega}_{1j} = \omega_{1j}$ for $k+1 \leq j \leq m$ and otherwise $\hat{\omega}_{1j} = 0$, $\hat{\omega}_{ij} = \omega_{ij}$ for all $i, j \geq 2$; $\hat{\lambda}_0 = \mu$, $\hat{\lambda}_1 = \nu$, $\hat{\lambda}_i = \lambda_i$ for all $i \geq 2$. This contradicts the fact that the rows of $R_{\text{cyl}}(\hat{G} - v_0v_1, \hat{p})$ are linearly independent. Hence $\dim \text{span}\{a, b\} = 2$.

Let \mathcal{P} be the plane which passes through $p(v_1)$ and whose normal belongs to $\text{span}\{a, b\}^\perp$. Since $\bar{p}(v_1) \in \text{span}\{a, b\}$, \mathcal{P} is not tangential to \mathcal{Y} at $p(v_1)$ and hence \mathcal{P} intersects \mathcal{Y} in a curve \mathcal{C} which passes through $p(v_1)$. Since $(\hat{G} - v_0v_1, \hat{p})$ is infinitesimally rigid, we have $\text{rank } R_{\text{cyl}}(\hat{G} - v_0v_1, q) = 3|V(\hat{G})| - 2$ for all q sufficiently close to \hat{p} on \mathcal{Y} . Hence we may choose a simple path $Q : [0, 1] \rightarrow \mathcal{C}$ such that $Q(0) = p(v_1)$ and such that the realisation (\hat{G}, q^t) with $q^t|_G = p$ and $q^t(v_0) = Q(t)$ is infinitesimally rigid for all $t \in [0, 1]$. Then $c^t = q^t(v_0) - p(v_1)$ satisfies $\text{span}\{a, b\} = \text{span}\{c^t, \bar{p}(v_1)\}$ for all $t \in (0, 1]$.

Let R_t be obtained from $R_{\text{cyl}}(\hat{G}, \hat{p})$ by replacing the zero row indexed by v_0v_1 by the row with c^t and $-c^t$ in the v_0 and v_1 columns and zeros elsewhere. The choice of c^t tells us that, for all $t \in (0, 1]$, there exist unique scalars $\omega_{01}^t, \bar{\omega}_{01}^t, \lambda_0^t, \lambda_1^t$ such that

$$(4.2) \quad a + \lambda_0^t \bar{p}(v_1) + \omega_{01}^t c^t = 0$$

and

$$(4.3) \quad b + \lambda_1^t \bar{p}(v_1) - \bar{\omega}_{01}^t c^t = 0.$$

Since (ω, λ) is an equilibrium stress for (G, p) we have

$$(4.4) \quad a + b + \lambda_1 \bar{p}(v_1) = 0.$$

Hence $(\lambda_0^t + \lambda_1^t) \bar{p}(v_1) + (\omega_{01}^t - \bar{\omega}_{01}^t) c^t = \lambda_1 \bar{p}(v_1)$. It follows that $\lambda_1 = \lambda_0^t + \lambda_1^t$ and $\omega_{01}^t = \bar{\omega}_{01}^t$.

We can extend $\omega_{01}^t, \lambda_0^t, \lambda_1^t$ to a vector $(\omega^t, \lambda^t) \in \text{coker } R_t$ as follows: $\omega_{0j}^t = \omega_{1j}$ for $2 \leq j \leq k$ and $\omega_{0j}^t = 0$ for $j > k$, $\omega_{1j}^t = \omega_{1j}$ for $k+1 \leq j \leq m$ and $\omega_{1j}^t = 0$ for $2 \leq j \leq k$ or $j > m$, $\omega_{ij}^t = \omega_{ij}$ for $i, j \geq 2$; $\lambda_i^t = \lambda_i$ for $i \geq 2$. Since $\text{rank } R_{\text{cyl}}(\hat{G} - v_0v_1, \hat{p}) = 3|V(\hat{G})| - 2$, the rows of R_t indexed by $E(\hat{G}) - v_0v_1$ are linearly independent. The fact that (ω^t, λ^t) is a non-zero vector in $\text{coker } R_t$ now gives $\omega_{01}^t \neq 0$.

The matrix $\Omega(\omega^t)$ defined by (ω^t, λ^t) has the form

$$\Omega(\omega^t) = \begin{bmatrix} \omega_{01}^t + \sum_{j \geq 2} \omega_{0j}^t & -\omega_{01}^t & -\omega_{02}^t & \cdots \\ -\omega_{01}^t & \omega_{01}^t + \sum_{j \geq 2} \omega_{1j}^t & -\omega_{12}^t & \cdots \\ -\omega_{02}^t & -\omega_{12}^t & \sum_{j \geq 2} \omega_{2j}^t & \cdots \\ \vdots & \vdots & \vdots & \vdots \end{bmatrix}.$$

Let $X = \begin{bmatrix} Y & 0 \\ 0 & I \end{bmatrix}$ be the $(n+1) \times (n+1)$ block matrix with $Y = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}$. We may use the fact that $\omega_{1j}^t + \omega_{0j}^t = \omega_{1j}$ for all $2 \leq j \leq k$ to obtain

$$X^T \Omega(\omega^t) X = \begin{bmatrix} \omega_{01}^t + \sum_{j \geq 2} \omega_{0j}^t & \sum_{j \geq 2} \omega_{0j}^t & -\omega_{02}^t & \cdots \\ \sum_{j \geq 2} \omega_{0j}^t & \sum_{j \geq 2} \omega_{1j} & -\omega_{12} & \cdots \\ -\omega_{02}^t & -\omega_{12} & \sum_{j \geq 2} \omega_{2j} & \cdots \\ \vdots & \vdots & \vdots & \vdots \end{bmatrix} = \begin{bmatrix} A_t & B \\ B^T & \Omega(\omega) \end{bmatrix}$$

where $A_t = (\omega_{01}^t + \sum_{j \geq 2} \omega_{0j}^t)$ and $B = (\sum_{j \geq 2} \omega_{0j}^t, -\omega_{02}^t, \dots, -\omega_{0n}^t)$. (Note that the entries in B do not change with t .) We will show that the entry in A_t is unbounded as $t \rightarrow 0$.

Let $\bar{p}(v_1)^\perp$ be a unit vector in $\text{span}\{a, b\}$ which is orthogonal to $\bar{p}(v_1)$. Then Equation (4.2) gives $a \cdot \bar{p}(v_1)^\perp + \omega_{01}^t c^t \cdot \bar{p}(v_1)^\perp = 0$. Since c^t and $\bar{p}(v_1)$ are linearly independent, we have $c^t \cdot \bar{p}(v_1)^\perp \neq 0$ and

$$(4.5) \quad \omega_{01}^t = -\frac{a \cdot \bar{p}(v_1)^\perp}{c^t \cdot \bar{p}(v_1)^\perp}.$$

As $t \rightarrow 0$, $q(v_0)$ approaches $p(v_1)$ on \mathcal{C} , and hence c^t will approach $(0, 0, 0)$. Hence $|\omega_{01}^t|$ will become arbitrarily large. On the other hand $|\omega_{0j}^t| \leq |\omega_{1j}|$ for all $j \geq 2$ and all $t > 0$. Hence $\lim_{t \rightarrow 0} |\omega_{01}^t + \sum_{j \geq 2} \omega_{0j}^t| = \infty$.

For all $t \in (0, 1]$, let $(\tilde{\omega}^t, \tilde{\lambda}^t)$ be the unique equilibrium stress for (\hat{G}, q^t) with $\tilde{\omega}_{01}^t = \omega_{01}^t$. Then $(\tilde{\omega}^t, \tilde{\lambda}^t)$ is a continuous function of t for $t \in (0, 1]$, and $\lim_{t \rightarrow 0} (\tilde{\omega}^t - \omega^t, \tilde{\lambda}^t - \lambda^t) = (0, 0)$.

It follows that $X^T \Omega(\tilde{\omega}^t) X = \begin{bmatrix} \tilde{A}_t & B_t \\ B_t^T & D_t \end{bmatrix}$ where B_t, D_t converge to B and $\Omega(\omega)$ respectively as $t \rightarrow 0$, and $\tilde{A}_t = (a_t)$ with $a_t = \omega_{01}^t + \sum_{j \geq 2} \tilde{\omega}_{0j}^t$ so $\lim_{t \rightarrow 0} |a_t| = \infty$. We can now use Lemma 4.3 to deduce that

$$\text{rank} \Omega(\tilde{\omega}^t) \geq 1 + \text{rank} \Omega(\omega)$$

when t is sufficiently close to zero.

We next consider the matrix $\Omega(\omega^t) + \Lambda(\lambda^t)$. We have

$$X^T (\Omega(\omega^t) + \Lambda(\lambda^t)) X = X^T \Omega(\omega^t) X + X^T \Lambda(\lambda^t) X = \begin{bmatrix} \bar{A}_t & B_t \\ B_t^T & \Omega(\omega) + \Lambda(\lambda) \end{bmatrix}$$

where $\bar{A}_t = A_t + (\lambda_0^t)$ and $B_t = B + (\lambda_0^t, 0, 0, \dots, 0)$. We will show that λ_0^t converges to a finite limit as $t \rightarrow \infty$. We can then use a similar argument to that in the previous paragraph to deduce that

$$\text{rank}(\Omega(\tilde{\omega}^t) + \Lambda(\tilde{\lambda}^t)) \geq 1 + \text{rank}(\Omega(\omega) + \Lambda(\lambda))$$

for all t sufficiently close to zero.

Let a^\perp be a unit vector in $\text{span}\{a, b\}$ which is orthogonal to a . Then Equation (4.2) gives $\lambda_0^t \bar{p}(v_1) \cdot a^\perp + \omega_{01}^t c^t \cdot a^\perp = 0$. If $\bar{p}(v_1) \cdot a^\perp = 0$ then $\bar{p}(v_1)$ would be a scalar multiple of a and Equation (4.4) would give $\dim \text{span}\{a, b\} = 1$. Hence $\bar{p}(v_1) \cdot a^\perp \neq 0$ and we may use Equation (4.5) to deduce that

$$\lambda_0^t = -\frac{\omega_{01}^t c^t \cdot a^\perp}{\bar{p}(v_1) \cdot a^\perp} = \frac{a \cdot \bar{p}(v_1)^\perp}{c^t \cdot \bar{p}(v_1)^\perp} \frac{c^t \cdot a^\perp}{\bar{p}(v_1) \cdot a^\perp} = \frac{a \cdot \bar{p}(v_1)^\perp}{\bar{p}(v_1) \cdot a^\perp} \frac{c^t \cdot a^\perp}{c^t \cdot \bar{p}(v_1)^\perp} = \frac{a \cdot \bar{p}(v_1)^\perp}{\bar{p}(v_1) \cdot a^\perp} \frac{\hat{c}^t \cdot a^\perp}{\hat{c}^t \cdot \bar{p}(v_1)^\perp}$$

where \hat{c}^t is a unit vector in the direction of c^t . As $t \rightarrow 0$, $q(v_0)$ approaches $p(v_1)$ on \mathcal{C} , and \hat{c}^t will approach $\bar{p}(v_1)^\perp$. Hence λ_0^t will become arbitrarily close to $\frac{a \cdot \bar{p}(v_1)^\perp \cdot \bar{p}(v_1)^\perp \cdot a^\perp}{\bar{p}(v_1) \cdot a^\perp}$ (since $\bar{p}(v_1)^\perp$ has unit length).

As noted above, the fact that λ_0^t converges to a finite limit as $t \rightarrow \infty$ implies that $\text{rank}(\Omega(\tilde{\omega}^t) + \Lambda(\tilde{\lambda}^t)) \geq 1 + \text{rank}(\Omega(\omega) + \Lambda(\lambda))$ for all t sufficiently close to 0. This in turn implies that $(\tilde{\omega}^t, \tilde{\lambda}^t)$ is an equilibrium stress for (\hat{G}, q^t) on \mathcal{Y} with

$$\text{rank } \Omega_{\text{cyl}}(\tilde{\omega}^t, \tilde{\lambda}^t) \geq \text{rank } \Omega_{\text{cyl}}(\omega, \lambda) + 3 = 3|V(\hat{G})| - 6$$

for all t sufficiently close to 0. The fact that equality holds follows since $3|V(\hat{G})| - 6$ is an upper bound on the rank of a stress matrix for any realisation of \hat{G} on \mathcal{Y} . \square

Combining the results thus far we have the following key result.

Theorem 4.7. *Let (G, p) be a generic realisation of an $\mathcal{M}_{2,2}^*$ -circuit on \mathcal{Y} . Then (G, p) has a maximum rank equilibrium stress.*

Proof. We apply induction on $|V|$. We give specific infinitesimally rigid realisations of K_5^- and H_1 on \mathcal{Y} which have maximum rank equilibrium stresses in the Appendix. Theorem 1.1 and Lemma 4.1 now imply that every generic realisation p of K_5^- or H_1 is infinitesimally rigid on \mathcal{Y} and has a maximum rank equilibrium stress on \mathcal{Y} . By Theorem 2.1 any $\mathcal{M}_{2,2}^*$ -circuit G can be formed from K_5^- or H_1 by K_4^- -extensions and generalised vertex splits. The result now follows from Theorems 4.2 and 4.6 and Lemma 4.1. \square

5. GLOBALLY RIGID FRAMEWORKS

In this section we will prove our main results, Theorems 1.2 and 1.3.

We say that a framework (G, p) on \mathcal{Y} is *quasi-generic* if it is congruent to a generic framework on \mathcal{Y} . The framework (G, p) is said to be in *standard position* on \mathcal{Y} if $p(v_1) = (0, 1, 0)$.

Lemma 5.1. *Let (G, p) and (G, p') be equivalent quasi-generic frameworks in standard position on \mathcal{Y} . Let $p(v_i) = (x_i, y_i, z_i)$ and $p'(v_i) = (x'_i, y'_i, z'_i)$ for each $v_i \in V$. Suppose that (G, p) has a maximum rank equilibrium stress (ω, λ) . Then $(z'_1, z'_2, \dots, z'_n) = b(z_1, z_2, \dots, z_n)$ for some $b \in \mathbb{R}$.*

Proof. By Theorem 1.6, (ω, λ') is an equilibrium stress for (G, p') for some $\lambda' \in \mathbb{R}$. Let

$$\Omega_{\text{cyl}}(\omega, \lambda) = \begin{bmatrix} \Omega + \Lambda & 0 & 0 \\ 0 & \Omega + \Lambda & 0 \\ 0 & 0 & \Omega \end{bmatrix} \quad \text{and} \quad \Omega_{\text{cyl}}(\omega, \lambda') = \begin{bmatrix} \Omega + \Lambda' & 0 & 0 \\ 0 & \Omega + \Lambda' & 0 \\ 0 & 0 & \Omega \end{bmatrix}$$

be the stress matrices corresponding to (ω, λ) and (ω, λ') , respectively. Observe that the submatrix Ω is the same in both stress matrices. We saw in Section 3 that $\text{rank } \Omega = n - 2$ and that $\{(z_1, z_2, \dots, z_n), (1, 1, \dots, 1)\}$ is a basis for $\text{coker } \Omega$. Since $(z'_1, z'_2, \dots, z'_n) \in \text{coker } \Omega$ we have $(z'_1, z'_2, \dots, z'_n) = b(z_1, z_2, \dots, z_n) + c(1, 1, \dots, 1)$ for some $b, c \in \mathbb{R}$. The hypothesis that $z_1 = z'_1 = 0$, now gives $c = 0$. \square

Our next result tells us that equivalent quasi-generic frameworks which satisfy the conclusion of Lemma 5.1 are in fact congruent. Its proof requires the introduction of a new concept. We will delay this to Section 7 and instead show how the result can be used to prove Theorem 1.3.

Lemma 5.2. *Let G be a 2-connected graph with at least $n + 1$ edges, and (G, p) and (G, p') be equivalent quasi-generic frameworks in standard position on \mathcal{Y} . Suppose that $p(v_i) = (x_i, y_i, z_i)$ and $p'(v_i) = (x'_i, y'_i, z'_i)$ for all $v_i \in V$, and that $(z'_1, z'_2, \dots, z'_n) = b(z_1, z_2, \dots, z_n)$ for some $b \in \mathbb{R}$. Then (G, p') is congruent to (G, p) .*

Proof of Theorem 1.3. Suppose that (G, p') is equivalent to (G, p) . We may assume that (G, p) is in fact quasi-generic and that (G, p) and (G, p') are both in standard position on \mathcal{Y} . Theorem 4.7 implies that (G, p) has a maximum rank equilibrium stress. Lemmas 5.1 and 5.2 now imply that (G, p) and (G, q) are congruent. \square

We will need some further concepts and results from matroid theory to deduce Theorem 1.2 from Theorem 1.3. A matroid is *connected* if any pair of edges is contained in a common circuit.

Lemma 5.3. [18, Theorem 5.4] *Suppose that G is a graph. Then $\mathcal{M}_{2,2}^*(G)$ is connected if and only if G is 2-connected and redundantly rigid on \mathcal{Y} .*

Let $\mathcal{M} = (E, \mathcal{J})$ be a matroid and let C_1, C_2, \dots, C_t be a non-empty sequence of circuits of \mathcal{M} . Let $D_j = C_1 \cup C_2 \cup \dots \cup C_j$ for $1 \leq j \leq t$ and suppose $D_t = E$. We say that C_1, C_2, \dots, C_t is an *ear decomposition* of \mathcal{M} if for all $2 \leq i \leq t$ the following properties hold:

- (1) $C_i \cap D_{i-1} \neq \emptyset$;
- (2) $C_i - D_{i-1} \neq \emptyset$;
- (3) no circuit C'_i satisfying (1) and (2) has $C'_i - D_{i-1}$ properly contained in $C_i - D_{i-1}$.

Lemma 5.4 ([7]). *A matroid is connected if and only if it has an ear decomposition.*

We also need a ‘glueing’ lemma for combining globally rigid frameworks. Its proof uses the following result.

Lemma 5.5. [13, Lemma 14] *Let G be a graph with at least five vertices and (G, p) and (G, q) be congruent generic realisations of G on the cylinder \mathcal{Y} . Then $\iota \circ p = q$ for some isometry ι of \mathcal{Y} .*

Note that the isometries of \mathcal{Y} are translations along and rotations about the z -axis, as well as reflections in any plane containing or orthogonal to the z -axis. Hence Lemma 5.5 implies that, if (G, p) and (G, q) are congruent generic frameworks on \mathcal{Y} with at least five vertices satisfying $p(v_1) = q(v_1)$ and $p(v_2) = q(v_2)$ for two distinct vertices v_1, v_2 of G , then $p = q$.

Lemma 5.6. *Let G_1 and G_2 be graphs on at least five vertices and with at least two vertices in common. Let (G, p) be a generic realisation of $G = G_1 \cup G_2$ on \mathcal{Y} and let $p_i = p|_{G_i}$. Suppose that (G_i, p_i) is globally rigid on \mathcal{Y} for $i = 1, 2$. Then (G, p) is globally rigid on \mathcal{Y} .*

Proof. Choose $u, v \in V(G_1) \cap V(G_2)$ and let (G, q) be an equivalent framework to (G, p) on \mathcal{Y} . By applying a suitable isometry of \mathcal{Y} to q we may assume that $p(u) = q(u)$. Since (G_1, p_1) is globally rigid on \mathcal{Y} , Lemma 5.5 tells us that $q|_{G_1} = \iota \circ p_1$ for some isometry ι of \mathcal{Y} . In particular $q(v) = (\iota \circ p)(v)$. Since (G_2, p_2) is globally rigid on \mathcal{Y} , Lemma 5.5 implies that there is a unique equivalent realisation of G_2 on \mathcal{Y} which maps u to $p(u)$ and v to $(\iota \circ p)(v)$. Since both $(G_2, q|_{G_2})$ and $(G_2, \iota \circ p_2)$ have this property, $q|_{G_2} = \iota \circ p_2$. Hence $q = \iota \circ p$ and (G, p) is congruent to (G, q) . \square

Proof of Theorem 1.2. Necessity follows from [13]. We prove sufficiency by induction on $|E|$. Since G is 2-connected and redundantly rigid, $\mathcal{M}_{2,2}^*(G)$ is connected by Lemma 5.3, so has an ear decomposition by Lemma 5.4. Let H_1, H_2, \dots, H_t be the $\mathcal{M}_{2,2}^*$ -circuits induced

by an ear decomposition of $\mathcal{M}_{2,2}^*(G)$ and $p_i = p|_{H_i}$. Then (H_i, p_i) is globally rigid on \mathcal{Y} for all $1 \leq i \leq t$ by Theorem 1.3. Since $|V(H_i)| \geq 5$ and $|(\cup_{j=1}^i V(H_j)) \cap V(H_{i+1})| \geq 2$ for all $1 \leq i < t$, we may deduce that (G, p) is globally rigid on \mathcal{Y} by repeated applications of Lemma 5.6. \square

Theorem 1.2 implies that global rigidity on the cylinder is a generic property. That is, if G is a graph and p is generic on \mathcal{Y} , then (G, p) is globally rigid on \mathcal{Y} if and only if (G, q) is globally rigid on \mathcal{Y} for all generic q . It also gives a polynomial time algorithm for checking generic global rigidity on the cylinder since 2-connectivity [23] and redundant rigidity [3, 16] can both be checked efficiently.

6. CYLINDRICAL FRAMEWORKS WITH ONE FREE VERTEX

In this section we characterise rigidity for a generic framework on the cylinder when one designated vertex is allowed to move off the cylinder. This is used to prove Lemma 4.5.

Let (G, p) be a generic realisation of an $\mathcal{M}_{2,2}^*$ -circuit on \mathcal{Y} , and (ω, λ) be a non-zero stress for (G, p) . If $\lambda_i = 0$ for some vertex v_i then $(\omega, \lambda|_{V-v_i})$ would be in the cokernel of the matrix $R_{v_i}(G, p)$ obtained from $R_{\text{cyl}}(G, p)$ by deleting the row indexed by v_i . Hence the rows of $R_{v_i}(G, p)$ would be dependent. Thus it will suffice to show that this cannot occur i.e. $\text{rank } R_v(G, p) = 3|V| - 2$ for all $v \in V$.

We can view $R_v(G, p)$ as a rigidity matrix for the framework (G, p) when v is free to move arbitrarily in \mathbb{R}^3 and all other vertices are constrained to stay on \mathcal{Y} . We will refer to such (infinitesimal) motions as *v-free (infinitesimal) motions* and say that (G, p) is *v-free (infinitesimally) rigid* if its only *v-free (infinitesimal) motions* are induced by the continuous isometries of \mathcal{Y} . Thus (G, p) is *v-free infinitesimally rigid* if and only if $\text{rank } R_v(G, p) = 3|V| - 2$ and Lemma 4.5 will follow immediately from our next result.

Theorem 6.1. *Let (G, p) be a generic realisation of an $\mathcal{M}_{2,2}^*$ -circuit on \mathcal{Y} and $v \in V$. Then (G, p) is v-free infinitesimally rigid.*

Note that ‘*v-free infinitesimal rigidity*’ is a generic property on \mathcal{Y} since it is characterised by the rank of $R_v(G, p)$. We will say that G is *v-free rigid on \mathcal{Y}* if some, or equivalently every, generic realisation of G on \mathcal{Y} is *v-free infinitesimally rigid*. Given graphs H and G , we say that H is a *0-extension* of G if H can be obtained from G by adding a vertex of degree two. We will need the following result in the proof of Theorem 6.1.

Lemma 6.2. *Let $G = (V, E)$ be a graph and $v \in V$.*

(a) *Suppose that G is a 0-extension of a graph H obtained by adding a new vertex v_0 and new edges v_0v_1 and v_0v_2 , $v \in V(H)$ and H is v-free rigid on \mathcal{Y} . Then G is v-free rigid on \mathcal{Y} . More precisely, if (G, p) is any framework on \mathcal{Y} such that $(H, p|_H)$ is v-free infinitesimal rigid and $p(v_0), p(v_1), p(v_2)$ are not collinear, then (G, p) is v-free infinitesimal rigid.*

(b) *Suppose that G is a 1-extension of a graph H , $v \in V(H)$ and H is v-free rigid on \mathcal{Y} . Then G is v-free rigid on \mathcal{Y} .*

(c) *Suppose that $G = H_1 \cup H_2$, $v \in V(H_1)$, H_1 is v-free rigid on \mathcal{Y} , H_2 is rigid on \mathcal{Y} and $H_1 \cap H_2 \neq \emptyset$. Then G is v-free rigid on \mathcal{Y} .*

(d) *Suppose that H is a subgraph of $G - v$ with at least four vertices, H is rigid on \mathcal{Y} , and G/H is v-free rigid on \mathcal{Y} . Then G is v-free rigid on \mathcal{Y} .*

Proof. (a) This follows since the definition of $R_v(G, p)$ implies that

$$\text{rank } R_v(G, p) = \text{rank } R_v(H, p|_H) + 3 = 3|V| - 2.$$

(b) Let $V = \{v_0, v_1, \dots, v_n\}$, let (G, p) be generic on \mathcal{Y} and suppose G is formed from H by deleting the edge v_1v_2 and then adding the vertex v_0 and edges v_0v_i for $1 \leq i \leq 3$. By symmetry between v_1 and v_2 , we may assume that $v_2 \neq v$.

Suppose that (G, p) is not v -free infinitesimally rigid on \mathcal{Y} . Then it follows that no framework (G, q) on \mathcal{Y} is v -free infinitesimally rigid. Let \mathcal{P} be the tangent plane to \mathcal{Y} at $p(v_2)$ and let a and b be orthogonal unit vectors in \mathcal{P} such that b is orthogonal to $p(v_2) - p(v_1)$. Consider a sequence of frameworks (G, q^k) on \mathcal{Y} , in which $q^k|_{V-v_0} = p|_{V-v_0}$ and $q^k(v_0)$ tends to $p(v_2)$ in the direction of a as $k \rightarrow \infty$. More precisely, the normalised vector $(p(v_2) - q^k(v_0))/\|p(v_2) - q^k(v_0)\|$ converges to a , as $k \rightarrow \infty$. Let $q = \lim_{k \rightarrow \infty} q^k$.

Since (G, q^k) is not v -free infinitesimally rigid, (G, q^k) has a unit norm infinitesimal motion m^k which has $m^k(v_2) = 0$ and hence is not an infinitesimal isometry of the cylinder. By the Bolzano-Weierstrass theorem there is a subsequence of the sequence m^k which converges to a vector, $m = (m_0, m_1, \dots, m_n)$ say. Relabeling if necessary, we may assume this holds for the original sequence m^k . Then m is a v -free infinitesimal motion of the framework (G, q) , for which $q(v_0) = p(v_2)$ and $q(v_i) = p(v_i)$ for all $1 \leq i \leq n$.

We claim that m is a v -free infinitesimal motion of $(G + v_1v_2, q)$. To see this, it will suffice to show that $m_2 - m_1$ is orthogonal to $p(v_2) - p(v_1)$. Since m is a v -free infinitesimal motion of (G, q) , $m_0, m_2 \in \mathcal{P}$ and, since $v_2v_0 \in E(G)$ and $(p(v_2) - q^k(v_0))/\|p(v_2) - q^k(v_0)\|$ converges to a , $m_2 - m_0$ is orthogonal to a . This tells us that $m_2 - m_0$ is a scalar multiple of b and hence that $m_2 - m_0$ is orthogonal to $p(v_2) - p(v_1)$. The fact that m is a v -free infinitesimal motion of (G, q) and $v_1v_0 \in E(G)$, also implies that $m_1 - m_0$ is orthogonal to $q(v_0) - q(v_1) = p(v_2) - p(v_1)$. Taking differences, we may deduce that $m_2 - m_1 = (m_2 - m_0) - (m_1 - m_0)$ is orthogonal to $p(v_2) - p(v_1)$, as desired.

Since the vectors m^k have unit norm and satisfy $m^k(v_2) = 0$, the vector m will also have unit norm and satisfy $m(v_2) = 0$. Hence $(G + v_1v_2, q)$ is not v -free infinitesimally rigid. This contradicts (a) since $(H, q|_H)$ is v -free infinitesimally rigid, $G + v_1v_2$ can be obtained from H by adding the edges v_0v_1, v_0v_3 (and v_0v_2), and $q(v_0) = p(v_2), q(v_1) = p(v_1), q(v_3) = p(v_3)$ are not collinear.

(c) Let (G, p) be a generic realisation of G on \mathcal{Y} and m be an infinitesimal motion of (G, p) . By applying an infinitesimal isometry of \mathcal{Y} to (G, p) we may assume that $m(v_0) = 0$ for some $v_0 \in V(H_1 - v)$. The fact that H_1 is v -free rigid on \mathcal{Y} now implies that the restriction of m to $(H_1, p|_{H_1})$ is zero. In particular, $m(u) = 0$ for all $u \in V(H_1) \cap V(H_2)$. Since H_2 is rigid on \mathcal{Y} and $H_1 \cap H_2 \neq \emptyset$, the restriction of m to $(H_2, p|_{H_2})$ is zero. Hence $m = 0$ and G is v -free rigid on \mathcal{Y} .

(d) Let $V = \{v_1, v_2, \dots, v_n\}$, $V(H) = \{v_1, v_2, \dots, v_r\}$ and (G, p) be a generic framework on \mathcal{Y} . By reordering its rows, we can write $R_v(G, p)$ in the form

$$\begin{pmatrix} R_{\text{cyl}}(H, p|_H) & 0 \\ M_1(p) & M_2(p) \end{pmatrix}$$

where $M_2(p)$ is a matrix with $3(n - r)$ columns.

Suppose, for a contradiction, that G is not v -free rigid on \mathcal{Y} . Then there exists a vector $m \in \ker R_v(G, p)$ which is not an infinitesimal isometry of \mathcal{Y} . Since $(H, p|_H)$ is rigid on \mathcal{Y} we may suppose that $m = (0, \dots, 0, m_{r+1}, \dots, m_n)$. Consider the realisation (G, p') where $p' = (p(v_r), p(v_r), \dots, p(v_r), p(v_{r+1}), \dots, p(v_n))$. Let v_r^* be the vertex of G/H which corresponds to H and define the realisation $(G/H, p^*)$ by setting $p^*(v_r^*) = p(v_r)$ and $p^*(v_i) = p(v_i)$ for all $r < i \leq n$. Since p^* is generic, $(G/H, p^*)$ is v -free rigid on \mathcal{Y} by assumption.

Since the nonzero vector $(m_{r+1}, \dots, m_n) \in \ker M_2(p)$, $\text{rank } M_2(p) < 3(n - r)$. Since p is generic, we also have $\text{rank } M_2(p') \leq \text{rank } M_2(p) < 3(n - r)$ and hence there exists a nonzero

vector $m' \in \ker M_2(p')$. Therefore we have

$$(R_v(G/H, p^*)) \begin{pmatrix} 0 \\ m' \end{pmatrix} = \begin{pmatrix} \star & M_2(p') \\ \bar{p}(v_r) & 0 \end{pmatrix} \begin{pmatrix} 0 \\ m' \end{pmatrix} = 0.$$

Thus $(0, m')$ is a nonzero infinitesimal motion of $(G/H, p^*)$ which fixes v_r^* . This contradicts the v -free rigidity of $(G/H, p^*)$. \square

Proof of Theorem 6.1 (and Lemma 4.5). We use a similar inductive proof technique to that of Theorem 2.3. The base graphs in our induction are the graphs K_5^-, H_1, H_2 . We give an infinitesimally rigid realisation (G, p) of G on \mathcal{Y} with a *nowhere zero* stress for each $G \in \{K_5^-, H_1, H_2\}$ in the Appendix. This implies that the rank of $R(G, p)$ is $3|V| - 2$ and that its rank will remain unchanged if we delete any of its rows. In particular, we have $\text{rank } R_v(G, p) = 3|V| - 2$ and hence G is v -free rigid on \mathcal{Y} . Thus we may assume that $G \notin \{K_5^-, H_1, H_2\}$.

Suppose that G has an admissible node x . If $x \neq v$ then the graph obtained from G by performing an admissible 1-reduction at x is v -free rigid on \mathcal{Y} by induction and hence G is v -free rigid on \mathcal{Y} by Lemma 6.2(b). On the other hand, if $x = v$ then v has degree three. Since G is an $\mathcal{M}_{2,2}^*$ -circuit, $G - v$ is rigid on \mathcal{Y} . This implies that G is v -free rigid on \mathcal{Y} since v has three neighbours in $G - v$.

Hence we may assume that G has no admissible nodes. Theorem 2.2 now implies that G has a nontrivial 2-separation or 3-edge-separation. It follows that G has at least two distinct atoms. Thus we may choose a nontrivial separation (F_1, F_2) of G such that F_1 is an atom of G and $v \in V(F_2)$. Consider the following cases.

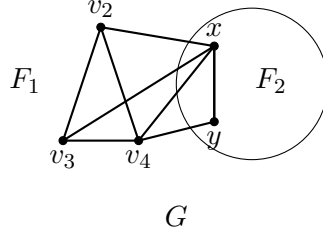
Case 1: (F_1, F_2) is a nontrivial 3-edge-separation.

Let G_i be obtained from G by contracting F_{3-i} to a single vertex z_i for each $i \in \{1, 2\}$. Then $G = G_1 *_3 G_2$ so G_1, G_2 are $\mathcal{M}_{2,2}^*$ -circuits on \mathcal{Y} by Theorem 1.4. Since G_1 is an $\mathcal{M}_{2,2}^*$ -circuit on \mathcal{Y} and $d(z_1) = 3$, $G_1 - z_1$ is rigid on \mathcal{Y} , and since G_2 is an $\mathcal{M}_{2,2}^*$ -circuit, G_2 is v -free rigid on \mathcal{Y} by induction. Lemma 6.2(d) now implies that G is v -free rigid on \mathcal{Y} .

Case 2: (F_1, F_2) is a nontrivial 2-vertex-separation, $V(F_1) \cap V(F_2) = \{x, y\}$ and $E(F_1) \cap E(F_2) = \{xy\}$.

Let G_i be obtained from F_i by adding two new vertices $\{w, z\}$ and five new edges $\{wx, wy, wz, xz, yz\}$. Then $G = G_1 *_2 G_2$ so G_1 and G_2 are $\mathcal{M}_{2,2}^*$ -circuits by Theorem 1.4. Since F_1 is an atom, G_1 has no nontrivial 2-vertex-separation or 3-edge-separation. If $G_1 \notin \{K_5^-, H_1, H_2\}$, then G_1 would have an admissible node $u \neq v$ by Theorem 2.2. It is straightforward to check that u would be an admissible node in G . Since G has no admissible nodes, we may deduce that $G_1 \in \{K_5^-, H_1, H_2\}$. We have $G_1 \neq K_5^-$ since G_1 is not 3-connected and $G_1 \neq H_1$ since the 2-separation (F_1, F_2) is nontrivial. Hence $G_1 = H_2$ and G is as shown in Figure 13.

Put $\{b, c\} = \{x, y\}$ with $b \neq v$. Let (G_2, p_2) be a generic realisation of G_2 on \mathcal{Y} . Since G_2 is an $\mathcal{M}_{2,2}^*$ -circuit, G_2 is v -free rigid by induction. Since $R_v(G_2, p)$ has $3|V(G_2)| - 2$ rows, $(G_2 - wz, p_2)$ will have a unique nonzero v -free infinitesimal motion which fixes b (up to scalar multiplication). Since w, z have degree two in $G_2 - wz$, $(F_2, p_2|_{F_2})$ also has a unique nonzero v -free infinitesimal motion m_2 which fixes b . The fact that G_2 is v -free rigid implies that $m_2(c) \neq 0$. Since G_1 is an $\mathcal{M}_{2,2}^*$ -circuit and $F_1 = G_1 - \{w, z\}$, we can choose a realisation (F_1, p_1) which is infinitesimally rigid on \mathcal{Y} and has $p_1(x) = p_2(x)$ and $p_1(y) = p_2(y)$. Let (G, p) be the realisation of G on \mathcal{Y} with $p|_{F_i} = p_i$. We will show that (G, p) is infinitesimally v -free rigid. We assume, for a contradiction, that (G, p) has a nonzero v -free infinitesimal motion m which keeps b fixed.

FIGURE 13. The graph G in the subcase $G_1 = H_2$ of Case 2.

Suppose that $v \notin \{x, y\}$. Then, since (F_1, p_1) is infinitesimally rigid on \mathcal{Y} , $m|_{F_1} = 0$. This contradicts the fact that we must have $m(c) \neq 0$, since $m|_{F_2}$ will be a nonzero scalar multiple of the motion m_2 defined in the previous paragraph. Hence $v = c$.

The fact that (G_1, p_1) is v -free rigid implies that (F_1, p_1) has a unique nonzero v -free infinitesimal motion m_1 which keeps b fixed (up to scalar multiplication). Thus $m(v) = \alpha m_1(v) = \beta m_2(v)$ for some $\alpha, \beta \in \mathbb{R}$. We will show that, as we vary the positions of the vertices of $F_1 - \{x, y\}$ in (G, p) , the direction of $m_1(v)$ changes. Since the direction of $m_2(v)$ remains fixed this will imply that we can choose p_1 such that (G, p) is infinitesimally v -free rigid.

When $v = y$ this follows since the K_4 subgraph of F_1 attached to x is rigid on \mathcal{Y} and hence $m_1|_{F_1 - y} = 0$. Thus $m_1(y)$ is normal to the plane through $p_1(x), p_1(v_4), p_1(y)$. Clearly the direction of $m_1(y)$ will change as we vary $p_1(v_4)$. Hence we may suppose that $v = x$. In this case we first apply an isometry of \mathcal{Y} to (G, p) so that $p(y) = (0, 1, 0)$ and then put $p(x) = (x_1, y_1, z_1)$, $p(v_2) = (1, 0, -1)$, $p(v_3) = (-1, 0, -1/3)$ and $p(v_4) = (1/\sqrt{2}, 1/\sqrt{2}, z_4)$. We may scale m_1 so that $m_1(x) = (1, u_5, u_6)$ and then solve $R_v(F_1, p_1)m_1 = 0$ to deduce that $u_5 = -\frac{(1+\sqrt{2})(\sqrt{2}x_1y_1z_4^2 - x_1z_1z_4 + z_1^2 - z_4z_1)}{z_4(2y_1z_4 + z_1 - 2z_4 + \sqrt{2}y_1z_4 - \sqrt{2}z_4)y_1}$. Hence the direction of $m_1(x)$ will change as we vary z_4 .

Case 3: (F_1, F_2) is a nontrivial 2-vertex-separation, $V(F_1) \cap V(F_2) = \{x, y\}$ and $xy \notin E$.

We have

$$|E(F_1)| + |E(F_2)| = |E(G)| = 2|V(G)| - 1 = 2|V(F_1)| + 2|V(F_2)| - 5$$

and $|E(F_i)| \leq 2|V(F_i)| - 2$ for each $i = 1, 2$ so $2|V(F_i)| - 3 \leq |E(F_i)| \leq 2|V(F_i)| - 2$. Consider the following subcases.

Subcase 3.1: $|E(F_1)| = 2|V(F_1)| - 3$. Let G_1 be obtained from F_1 by adding two new vertices $\{w, z\}$ and six new edges $\{wx, wy, wz, xy, xz, yz\}$, and $G_2 = F_2 + xy$. Then $G = G_1 *_1 G_2$ so G_1 and G_2 are $\mathcal{M}_{2,2}^*$ -circuits by Theorem 1.4. Since F_1 is an atom, G_1 has no nontrivial 2-vertex-separation or 3-edge-separation. If $G_1 \notin \{K_5^-, H_1, H_2\}$, then G_1 would have an admissible node $u \neq v$ by Theorem 2.2. It is straightforward to check that u would be an admissible node in G . Since G has no admissible nodes we may deduce that $G_1 \in \{K_5^-, H_1, H_2\}$. We have $G_1 \neq K_5^-$ since G_1 is not 3-connected. Hence $G_1 \in \{H_1, H_2\}$.

Suppose $G_1 = H_1$. Then $F_1 = K_4^-$ and G is a K_4^- -extension of G_2 . Let (G_2, p_2) be a generic realisation of G_2 on \mathcal{Y} . Then (G_2, p_2) is v -free infinitesimally rigid by induction. Let $(G + xy, p)$ be a realisation of $G + xy$ on \mathcal{Y} such that $p|_{F_2} = p_2$ and such that the vertices of $p|_{F_1}$ are coplanar and are in general position on this plane. Then $(G + xy, p)$ is v -free infinitesimally rigid since it can be obtained from (G_2, p_2) by two 0-extensions and an edge

addition. In addition the rows of $R_v(G + xy)$ indexed by $E(F_1) + xy$ are minimally linearly dependent since the vertices of $p|_{F_1}$ are coplanar and are in general position on this plane. Hence we may delete the edge xy from $(G + xy, p)$ without destroying its v -free infinitesimal rigidity. Thus we may assume that $G_1 \neq H_1$.

Suppose $G_1 = H_2$. Then G is as shown in Figure 14(i). If $v \neq x$ then we may apply Lemma 6.2(d) with $H = F_1 - y$ to deduce that G is v -free rigid on \mathcal{Y} . Hence $v = x$. Let (G_2, p_2) be a generic realisation of G_2 on \mathcal{Y} . Since $(G_2, p|_{G_2})$ is v -free rigid by induction, $(F_2, p|_{F_2})$ has a unique nonzero infinitesimal motion m_2 which keeps y fixed (up to scalar multiplication), and $m_2(x) \neq 0$.

Since G_1 is v -free rigid, we can choose a realisation (F_1, p_1) such that p_1 agrees with p_2 at x, y and such that the space, Z_1 , of infinitesimal motions of (F_1, p_1) which keep y fixed, is 2-dimensional. Let (G, p) be the framework with $p|_{F_i} = p_i$. We will show that we can choose p_1 such that the infinitesimal velocities $m_1(x)$ for $m_1 \in Z_1$ span a 2-dimensional subspace \mathcal{P} of \mathbb{R}^3 , and that, as we vary the positions of the vertices of $F_1 - \{x, y\}$, the normals to \mathcal{P} span \mathbb{R}^3 . This will imply that we can choose p_1 so that no nonzero infinitesimal velocity $m_1(x)$ for $m_1 \in Z_1$, has the same direction as $m_2(x)$, and hence that (G, p) is infinitesimally rigid.

We again apply an isometry of \mathcal{Y} to (G, p) so that $p(y) = (0, 1, 0)$, and choose $p(x)$ and $p(v_i)$, $2 \leq i \leq 4$, to be as defined in the last paragraph of Case 2. Then (F_1, p_1) has two linearly independent infinitesimal motions m_1, \tilde{m}_1 with $m_1(y) = (0, 0, 0) = \tilde{m}_1(y)$, $m_1(x) = (1, 0, -\frac{-\sqrt{2}x_1z_4 + 2x_1y_1z_4 + \sqrt{2}z_1 - \sqrt{2}z_4}{2z_4y_1(z_1 - z_4)})$ and $\tilde{m}_1(x) = (0, 1, -\frac{-\sqrt{2} + 2y_1}{2(z_1 - z_4)})$. The normal vector to \mathcal{P} is $n_1 = m_1(x) \times \tilde{m}_1(x) = (\frac{-\sqrt{2}x_1z_4 + 2x_1y_1z_4 + \sqrt{2}z_1 - \sqrt{2}z_4}{2z_4y_1(z_1 - z_4)}, \frac{-\sqrt{2} + 2y_1}{2(z_1 - z_4)}, 1)$. We can obtain three linearly independent normal vectors by evaluating n_1 at each $z_4 \in \{1, 2, -1\}$. Hence the normals to \mathcal{P} span \mathbb{R}^3 as required.

Subcase 3.2: $|E(F_1)| = 2|V(F_1)| - 2$. Let G_2 be obtained from F_2 by adding two new vertices $\{w, z\}$ and six new edges $\{wx, wy, wz, xy, xz, yz\}$, and $G_1 = F_1 + xy$. Then $G = G_1 *_1 G_2$ so G_1 and G_2 are $\mathcal{M}_{2,2}^*$ -circuits by Theorem 1.4. Let (G_2, p_2) be a generic realisation of G_2 on \mathcal{Y} . Since G_1 is an $\mathcal{M}_{2,2}^*$ -circuit and $F_1 = G_1 - xy$, we can choose a realisation (G_1, p_1) such that (F_1, p_1) is infinitesimally rigid on \mathcal{Y} and p_1 agrees with p_2 on $\{x, y\}$. Let (G, p) be the framework with $p|_{F_i} = p_i$.

Suppose $v \notin \{x, y\}$. Since $v \in V(F_2)$ we have $v \notin V(F_1)$. Let m be an infinitesimal motion of (G, p) which keeps y fixed. Since (F_1, p_1) is infinitesimally rigid on \mathcal{Y} , $m|_{F_1} = 0$ and, in particular, $m(x) = 0$. It now follows that $m|_{F_2}$ can be extended to an infinitesimal motion m_2 of (G_2, p_2) by putting $m_2(w) = m_2(z) = 0$. Since (G_2, p_2) is infinitesimally v -free rigid by induction and $m_2(y) = 0$, we have $m_2 = 0$. Hence $m = 0$ and (G, p) is infinitesimally v -free rigid. Thus we may assume that $v \in \{x, y\}$.

Let $\{x, y\} = \{v, b\}$. Since G_2 is v -free rigid by induction, the space Z_2 of infinitesimal motions of $(F_2, p|_{F_2})$ which keep b fixed is 2-dimensional. Similarly, since G_1 is v -free rigid by induction, the space Z_1 of infinitesimal motions of $(F_1, p|_{F_1})$ which keep b fixed is 1-dimensional. Our strategy to show that (G, p) is v -free infinitesimally rigid is to adjust the positions of the vertices of $F_1 - \{x, y\}$ so that, for all nonzero $m_1 \in Z_1$ and $m_2 \in Z_2$, $m_1(v) \neq m_2(v)$. We will accomplish this by showing that the vectors $m_1(v)$ span \mathbb{R}^3 as $p_1|_{F_1 - \{x, y\}}$ varies.

Suppose v has degree three in G_1 and let b, c, d be the neighbours of v in G_1 . Since G_1 is an $\mathcal{M}_{2,2}^*$ -circuit, $G_1 - v$ is rigid on \mathcal{Y} and hence $m_1(u) = 0$ for all $u \in V(G_1) - v$ and all $m_1 \in Z_1$. It follows that $m_1(v)$ is normal to the plane through $p(v), p(c)$ and $p(d)$ for all $m_1 \in Z_1$. This implies that we can choose $p(c)$ and $p(d)$ to ensure that $m_1(v) \neq m_2(v)$ for

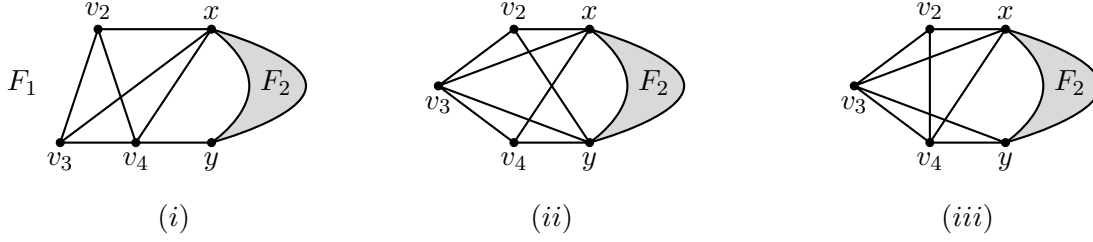


FIGURE 14. The graph G in the subcases $G_1 = H_2$ of Subcase 3.1 and $G_1 = K_5^-$ of Subcase 3.2.

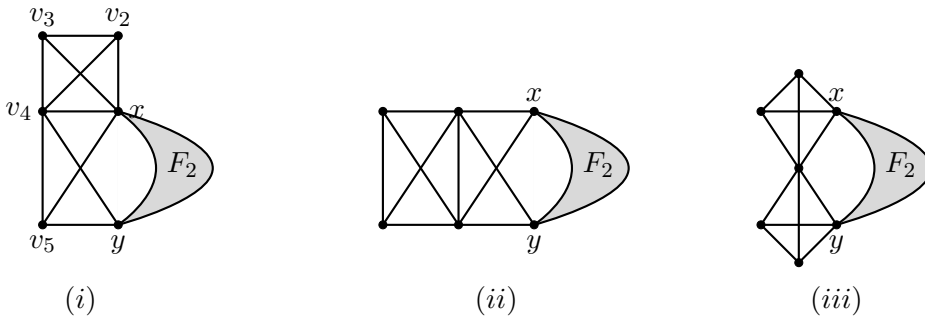


FIGURE 15. The graph G in the subcases $G_1 = H_1$ and $G_1 = H_2$ of Subcase 3.2.

all $m_1 \in Z_1$ with $m_1(v) \neq 0$ and all $m_2 \in Z_2$. Hence we may assume that v has degree at least four in G_1 .

Since F_1 is an atom, G_1 has no nontrivial 2-vertex-separation or 3-edge-separation. If $G_1 \notin \{K_5^-, H_1, H_2\}$, then G_1 has two admissible nodes by Theorem 2.2. Since v has degree at least four in G_1 , G_1 has at least one admissible node which is distinct from v, b , and this node is an admissible node of G . This contradicts the fact that G has no admissible nodes and hence we may assume that $G_1 \in \{K_5^-, H_1, H_2\}$.

Suppose $G_1 = K_5^-$. Then G has the structure of one of the graphs shown in Figure 14(ii) and (iii). If case (ii) occurs then we may assume by symmetry that $v = x$. On the other hand, if case (iii) occurs then we must also have $v = x$ since v has degree at least four in G_1 . In both cases we consider the realisation (G, p) where $p|_{F_1}$ is as defined in the last paragraph of Case 2, and determine the unique vector $m_1(v) = (1, u_5, u_6)$, for $m_1 \in Z_1$. We then show that the vectors $m_1(v)$ span \mathbb{R}^3 as the positions of the vertices of $F_1 - \{x, y\}$ vary. For simplicity in the expressions we also specify that $x_1 = 0, y_1 = -1$ and $z_1 = 1$.

When case (ii) occurs we obtain $u_5 = \frac{30z_4^2 - 5z_4\sqrt{2} - 12z_4 + 3\sqrt{2} + 2}{-z_4\sqrt{2} + 6z_4^2 - 3z_4}$ and $u_6 = \frac{-3z_4\sqrt{2} + 18z_4^2 - 3z_4 + 6\sqrt{2} + 4}{-z_4\sqrt{2} + 6z_4^2 - 3z_4}$ and when case (iii) occurs we obtain $u_5 = \frac{-\sqrt{2}z_4 + 6z_4^2 - 6z_4 + \sqrt{2}}{z_4(-\sqrt{2} + 6z_4 + 2)}$ and $u_6 = \frac{2(3z_4 - \sqrt{2})}{z_4(-\sqrt{2} + 6z_4 + 2)}$. In both cases, we can obtain three linearly independent infinitesimal velocities $m_1(v)$ by choosing $z_4 \in \{1, -1, 2\}$ in the above expressions for $(1, u_5, u_6)$.

Suppose $G_1 = H_1$. Since F_1 is an atom, G would have the structure of one of the graphs shown in Figure 15(i) and (ii). Since v has degree at least four in G_1 , case (ii) cannot occur and $v = x$ when case (i) occurs. Consider the realisation (G, p) where $p|_{F_1}$ is given by $p(y) = (0, 1, 0)$, $p(x) = (x_1, y_1, z_1)$, $p(v_2) = (-1, 0, -1/3)$, $p(v_3) = (1/\sqrt{2}, -1/\sqrt{2}, 1/3)$, $p(v_4) = (1, 0, -1)$ and $p(v_5) = (0, -1, z_5)$. For simplicity in the expressions we also specify

that $x_1 = -1/\sqrt{2}$, $y_1 = 1/\sqrt{2}$ and $z_1 = -2$. Then the unique vector $m_1(v) = (1, u_5, u_6)$ with $m_1 \in Z_1$ is given by $u_5 = \frac{(49\sqrt{2}+73)(17z_5-71+77\sqrt{2})}{17(31z_5+111+21\sqrt{2})}$ and $u_6 = \frac{(-1+4\sqrt{2})(3z_5+17+11\sqrt{2})}{31z_5+111+21\sqrt{2}}$. By putting $z_5 = 0, 1$ in the above expressions for $(1, u_5, u_6)$ we get the infinitesimal velocities:

$$m_1(v) = \left(1, \frac{(49\sqrt{2} + 73)(-71 + 77\sqrt{2})}{17(111 + 21\sqrt{2})}, \frac{(-1 + 4\sqrt{2})(17 + 11\sqrt{2})}{111 + 21\sqrt{2}} \right)$$

and

$$\tilde{m}_1(v) = \left(1, \frac{(49\sqrt{2} + 73)(-54 + 77\sqrt{2})}{17(142 + 21\sqrt{2})}, \frac{(-1 + 4\sqrt{2})(20 + 11\sqrt{2})}{142 + 21\sqrt{2}} \right).$$

We also consider the realisation (G, \hat{p}) where $\hat{p}|_{F_2} = p|_{F_2}$ and $\hat{p}|_{F_1 - \{x, y\}}$ is given by $\hat{p}(v_2) = (1/\sqrt{2}, -1/\sqrt{2}, -1/2)$, $\hat{p}(v_3) = (-1, 0, 1/2)$, $\hat{p}(v_4) = (1/\sqrt{2}, 1/\sqrt{2}, 1)$ and $\hat{p}(v_5) = (1, 0, z_5)$. For simplicity in the expressions we again specify that $x_1 = -1/\sqrt{2}$, $y_1 = 1/\sqrt{2}$ and $z_1 = -2$. Then the unique vector $\hat{m}_1(v) = (1, u_5, u_6)$ with $\hat{m}_1 \in Z_1$ is given by $u_5 = \frac{\sqrt{2}z_5 + 2z_5^2 - 9\sqrt{2} + 11z_5 - 12}{4\sqrt{2}z_5 + 2z_5^2 - 15\sqrt{2} + 17z_5 - 18}$ and $u_6 = -\frac{5z_5 - 12 + \sqrt{2}z_5 - 3\sqrt{2} + 2z_5^2}{4\sqrt{2}z_5 + 2z_5^2 - 15\sqrt{2} + 17z_5 - 18}$. By putting $z_5 = 0$ in the this expression for $\hat{m}_1(v)$ we get the infinitesimal velocity

$$\hat{m}_1(v) = \left(1, \frac{-12 - 9\sqrt{2}}{-18 - 15\sqrt{2}}, -\frac{-12 - 3\sqrt{2}}{-18 - 15\sqrt{2}} \right).$$

We can now check that $m_1(v)$, $\tilde{m}_1(v)$ and $\hat{m}_1(v)$ are linearly independent.

The only remaining alternative is $G_1 = H_2$. Since F_1 is an atom, G would have the structure of the graph shown in Figure 15(iii). By symmetry we may assume that $v = x$. We can now use Lemma 6.2(d), with H equal to the copy of K_4 attached at y , and induction to deduce that G is v -free rigid. \square

We close this section by using Theorem 6.1 to obtain a full characterisation of v -free rigidity on the cylinder.

Theorem 6.3. *Let $G = (V, E)$ be a graph and $v \in V$. Then G is v -free rigid on \mathcal{Y} if and only if G is rigid on \mathcal{Y} and v is contained in an $\mathcal{M}_{2,2}^*$ -circuit of G .*

Proof. Suppose G is infinitesimally v -free rigid and choose a minimally v -free rigid spanning subgraph $H = (V, E)$ of G . Then H is rigid on \mathcal{Y} and $|E| = 2|V| - 1$, so H contains a unique $\mathcal{M}_{2,2}^*$ -circuit C . If $v \notin V(C)$ then the rows of $R_v(H, p)$ indexed by C would be linearly dependent for any generic (H, p) and hence H would not be minimally v -free rigid. Hence $v \in V(C)$.

Conversely, suppose that $G = (V, E)$ is rigid on \mathcal{Y} and v is contained in an $\mathcal{M}_{2,2}^*$ -circuit C of G . We will show that G is infinitesimally v -free rigid by induction on $|V| + |E|$. Choose a spanning subgraph H of G such that H is rigid on \mathcal{Y} , $C \subseteq H$ and H has as few edges as possible. If $H \neq G$ then we are through by induction, so we may assume that $H = G$. Hence $|E| = 2|V| - 1$ and C is the unique $\mathcal{M}_{2,2}^*$ -circuit in G .

Let H_1 be a maximal v -free rigid subgraph of G with $C \subseteq H_1$. Note that H_1 exists since C itself is v -free rigid by Theorem 6.1. We may assume that $V(G - H_1) \neq \emptyset$. The facts that $|E| = 2|V| - 1$ and $|E(H_1)| = 2|V(H_1)| - 1$ now imply that there exists a vertex u of degree two or three in G which does not belong to H_1 . If u has degree two then we can apply induction to $G - u$ and then use Lemma 6.2(a) to deduce that G is v -free rigid. Hence we may assume that $d(u) = 3$ and let $N(u) = \{x, y, z\}$. Since G has a unique $\mathcal{M}_{2,2}^*$ -circuit, at most two of x, y, z are contained in $V(H_1)$.

Suppose u, x, y, z induces a K_4 in G . Let H_2 be a maximal subgraph of G such that $u, x, y, z \in V(H_2)$ and $|E(H_2)| = 2|V(H_2)| - 2$. Then H_2 is rigid on \mathcal{Y} by Theorem 1.2. If $H_1 \cap H_2 \neq \emptyset$ then $H_1 \cup H_2$ is v -free rigid by Lemma 6.2(c). This contradicts the maximality of H_1 and hence $H_1 \cap H_2 = \emptyset$. Since G/H_2 is v -free rigid on \mathcal{Y} by induction, we can now use Lemma 6.2(d) to deduce that G is v -free rigid on \mathcal{Y} .

Hence we may suppose that u, x, y, z does not induce a K_4 in G . If we can apply a 1-reduction at u to obtain a graph G' which has $|E(G')| = 2|V(G')| - 1$ and contains a unique $\mathcal{M}_{2,2}^*$ -circuit, then we can apply induction and Lemma 6.2(b) to deduce that G is v -free rigid on \mathcal{Y} . Hence we may assume that, for all $a, b \in N(u)$ with $\{a, b\} \not\subseteq V(C)$, either $ab \in E$ or there exists a set $X_{ab} \subseteq V - u$ which induces $2|X_{ab}| - 2$ edges in G and has $a, b \in X_{ab}$.

Suppose at most one of x, y, z is in $V(C)$. Since u, x, y, z does not induce a K_4 in G we may assume that X_{xy} exists. Since C is the unique $\mathcal{M}_{2,2}^*$ -circuit in G , $z \notin X_{xy}$. If $\{xz, yz\} \subset E$ then $G[X_{xy} \cup \{u, z\}]$ would contain an $\mathcal{M}_{2,2}^*$ -circuit distinct from C . Hence we may suppose that $xz \notin E$ and X_{xz} exists. Then $G[X_{xy} \cup X_{xz} \cup \{u\}]$ contains an $\mathcal{M}_{2,2}^*$ -circuit distinct from C .

We may now assume that $x, y \in V(C)$ and $z \notin V(C)$. If $\{xz, yz\} \subseteq E$ then $(C - e) \cup G[u, x, y, z]$ would contain an $\mathcal{M}_{2,2}^*$ -circuit distinct from C for any $e \in E(C)$. Hence we may assume that X_{xz} exists. Then $(C - e) \cup G[X_{xz} \cup \{u\}]$ again contradicts the uniqueness of C . \square

7. VERTICALLY RESTRICTED FRAMEWORKS

In this section we prove Lemma 5.2. We will assume throughout the section that all frameworks are in standard position on \mathcal{Y} , with respect to a fixed vertex v_1 . Given a framework (G, \hat{p}) on \mathcal{Y} , we are interested in the set of frameworks (G, p) on \mathcal{Y} which are equivalent to (G, \hat{p}) and satisfy $(z_1, z_2, \dots, z_n) = b(\hat{z}_1, \hat{z}_2, \dots, \hat{z}_n)$ for some $b \in \mathbb{R}$, where $\hat{p}(v_i) = (\hat{x}_i, \hat{y}_i, \hat{z}_i)$ and $p(v_i) = (x_i, y_i, z_i)$. We will say that such a framework (G, p) is *vertically restricted equivalent* to (G, \hat{p}) . Henceforth we will abbreviate ‘vertically restricted’ by the prefix VR. We say that two frameworks are *VR-congruent* if they are linked by a reflection in the xy -plane or the yz -plane. Thus VR-congruence implies VR-equivalence. We say that (G, \hat{p}) is *globally VR-rigid* if every VR-equivalent framework is VR-congruent to it. When $\hat{z}_2 \neq 0$, the set of all frameworks in standard position with respect to v_1 which are VR-equivalent to (G, \hat{p}) is given by the set of solutions to the following system of equations:

$$(7.1) \quad \|p(v_i) - p(v_j)\|^2 = c_{ij} \quad (v_i v_j \in E)$$

$$(7.2) \quad \frac{z_i}{z_2} = k_i \quad (v_i \in V \setminus \{v_1, v_2\})$$

$$(7.3) \quad x_i^2 + y_i^2 = 1 \quad (v_i \in V - v_1)$$

where $\hat{p}(v_1) = (0, 1, 0)$, $c_{ij} = \|\hat{p}(v_i) - \hat{p}(v_j)\|^2$ and $k_i = \frac{\hat{z}_i}{\hat{z}_2}$. We can differentiate these equations to obtain the following linear system for the unknowns $\dot{p}(v_i)$, $v_i \in V - v_1$:

$$(7.4) \quad (p(v_i) - p(v_j)) \cdot (\dot{p}(v_i) - \dot{p}(v_j)) = 0 \quad (v_i v_j \in E)$$

$$(7.5) \quad z_2 \dot{z}_i - z_i \dot{z}_2 = 0 \quad (v_i \in V \setminus \{v_1, v_2\})$$

$$(7.6) \quad x_i \dot{x}_i + y_i \dot{y}_i = 0 \quad (v_i \in V - v_1)$$

where $\dot{p}(v_1) = 0$. We say that (G, \hat{p}) is *infinitesimally VR-rigid* if the only solution to this linear system is the trivial solution $\dot{p}(v_i) = 0$ for all $v_i \in V - v_1$, or equivalently if the rank of the matrix of coefficients of the system is $3|V| - 3$. This matrix, the *VR-rigidity matrix*

of (G, p) , has $3|V| - 3$ columns and $|E| + 2|V| - 3$ rows. The rows corresponding to (7.4) have the form

$$(\dots \ 0 \ p(v_i) - p(v_j) \ 0 \ \dots \ 0 \ p(v_j) - p(v_i) \ 0 \ \dots)$$

when $v_i, v_j \neq v_1$, and

$$(\dots \ 0 \ 0 \ 0 \ \dots \ 0 \ p(v_j) - p(v_1) \ 0 \ \dots)$$

when $v_i = v_1$. The rows corresponding to (7.5) have the form

$$((0, 0, -z_i) \ 0 \ \dots \ 0 \ (0, 0, z_2) \ 0 \ \dots).$$

The rows corresponding to (7.6) have the form

$$(\dots \ 0 \ (x_i, y_i, 0) \ 0 \ \dots).$$

We say that the graph G is *VR-rigid* if (G, p) is infinitesimally VR-rigid for some, or equivalently every, quasi-generic framework (G, p) in standard position on \mathcal{Y} . The graph G is *minimally VR-rigid* if G is VR-rigid, and $G - e$ is not rigid for all edges e of G . We first characterise minimal VR-rigidity.

Theorem 7.1. *A graph G is minimally VR-rigid if and only if it is connected and contains exactly one cycle.*

Proof. It is straightforward to show that if G is minimally VR-rigid then G is connected and has $|V|$ edges. To prove the converse, we first observe that every connected graph with exactly one cycle can be constructed recursively from K_3 by subdividing edges and adding vertices of degree one. The structure of the VR-rigidity matrix implies that K_3 is minimally VR-rigid and that adding vertices of degree one preserves minimal rigidity. To see that subdivision preserves minimal rigidity we can apply a similar technique to the proof of Lemma 6.2(b). (The new proof is simpler since there is no free vertex and our restriction to frameworks in standard position allows no infinitesimal isometry of the cylinder.) \square

The remaining results of this section use the algebraic independence of coordinates of a quasi-generic framework in standard position on \mathcal{Y} . The proof technique is similar to that in [11, 13]. In particular we use the following analogue of the ‘ \mathcal{Y} -rigidity map’ from [13]. Given a graph $G = (V, E)$ with n vertices and m edges, we define three maps with domain $D = \{(x_2, y_2, z_2, \dots, x_n, y_n, z_n) \in \mathbb{R}^{3n-3} : z_2 \neq 0\}$. We first associate a framework (G, p) in standard position on \mathcal{Y} with each $\tilde{p} \in D$ by putting $p = (0, 1, 0, \tilde{p})$. The first map $f_G : D \rightarrow \mathbb{R}^m$ is defined by $f_G(\tilde{p}) = (\|e_1\|^2, \dots, \|e_m\|^2)$ where $\|e_i\|^2 = \|p(v_j) - p(v_k)\|^2$ for each $e_i = v_j v_k \in E$. The second map $h_G : D \rightarrow \mathbb{R}^{n-2}$ is given by $h_G(\tilde{p}) = (z_3/z_2, z_4/z_2, \dots, z_n/z_2)$. The third map $\theta_G : D \rightarrow \mathbb{R}^{n-1}$ is given by $\theta_G(\tilde{p}) = (x_2^2 + y_2^2, \dots, x_n^2 + y_n^2)$. The *VR-rigidity map* $F_G^{\text{VR}} : D \rightarrow \mathbb{R}^{m+2n-3}$ is defined by $F_G^{\text{VR}} = (f_G, h_G, \theta_G)$. We showed in [13, Lemma 8] that a framework (G, p) is quasi-generic and in standard position on \mathcal{Y} if and only if $\text{td}[\mathbb{Q}(p) : \mathbb{Q}] = 2n - 2$. We may use a similar proof technique to [13, Lemma 9] to show further, that if G is VR-rigid and (G, p) is quasi-generic and in standard position on \mathcal{Y} , then $\overline{\mathbb{Q}(p)} = \overline{\mathbb{Q}(f_G(p), h_G(p))}$. This gives

Lemma 7.2. *Suppose (G, p) and (G, q) are two VR-equivalent frameworks in standard position on \mathcal{Y} and that (G, p) is infinitesimally VR-rigid and quasi-generic. Then $\overline{\mathbb{Q}(p)} = \overline{\mathbb{Q}(f_G(p), h_G(p))} = \overline{\mathbb{Q}(q)}$ and $\text{td}[\mathbb{Q}(p) : \mathbb{Q}] = 2|V| - 2$.*

We say that a graph G is *redundantly VR-rigid* on \mathcal{Y} if $G - e$ is VR-rigid on \mathcal{Y} for all edges e of G . Our next result shows that 2-connectivity and redundant rigidity are necessary conditions for generic global VR-rigidity.

Lemma 7.3. *Let (G, p) be a quasi-generic, globally VR-rigid framework in standard position on \mathcal{Y} . Then G is 2-connected and redundantly VR-rigid on \mathcal{Y} .*

Proof. Suppose G is not 2-connected. Then we have $G = G_1 \cup G_2$ for subgraphs $G_i = (V_i, E_i)$ with $|V_1 \cap V_2| \leq 1$ and $v_1 \in V_1$. Let $p_1 = p|_{V_1}$ and $p_2 = p|_{V_2}$. Let (G, q) be obtained from (G, p) by reflecting (G_2, p_2) in a plane which contains $p(V_1 \cap V_2)$ and also contains the z -axis. Then (G, q) is a framework in standard position on \mathcal{Y} and is VR-equivalent but not VR-congruent to (G, p) . Hence G is 2-connected.

Suppose $G - e$ is not VR-rigid for some $e = v_i v_j \in E$. Since p is quasi-generic, we can use a similar argument to that given in [13] to deduce that the configuration space of all frameworks in standard position on \mathcal{Y} which are VR-equivalent to $(G - e, p)$ is a 1-dimensional manifold. Let \mathcal{C} be the component of this manifold which contains p . Then \mathcal{C} is bounded since $G - e$ is connected. The fact that \mathcal{C} is closed now implies that \mathcal{C} is diffeomorphic to a circle. We can now use a similar argument to that given in [13, Theorem 9] to find a $q \in \mathcal{C}$ such that (G, q) is VR-equivalent, but not VR-congruent, to (G, p) . \square

Theorem 7.4. *A quasi-generic framework (G, p) in standard position on \mathcal{Y} is globally VR-rigid if and only if $G = (V, E)$ is 2-connected and $|E| \geq |V| + 1$.*

Proof. Necessity follows from Lemma 7.3. To prove the converse, we assume that (G, q) is a quasi-generic framework in standard position on \mathcal{Y} which is VR-equivalent to (G, p) . Let $p(v_i) = (x_i, y_i, z_i)$ and $q(v_i) = (a_i, b_i, c_i)$ for all $v_i \in V$.

We first consider the case when G contains a vertex v_0 of degree two which is distinct from v_1 . Let v_j, v_k be the neighbours of v_0 in G . By symmetry we may assume that $v_j \neq v_k$. We have the following constraints on the positions of v_0, v_j, v_k :

$$(7.7) \quad (x_0 - x_j)^2 + (y_0 - y_j)^2 + (z_0 - z_j)^2 - (a_0 - a_j)^2 - (b_0 - b_j)^2 - (c_0 - c_j)^2 = 0$$

$$(7.8) \quad (x_0 - x_k)^2 + (y_0 - y_k)^2 + (z_0 - z_k)^2 - (a_0 - a_k)^2 - (b_0 - b_k)^2 - (c_0 - c_k)^2 = 0$$

$$(7.9) \quad -c_0 z_j + c_j z_0 = 0$$

and

$$(7.10) \quad x_i^2 + y_i^2 = 1 = a_i^2 + b_i^2 \text{ for all } i \in \{0, j, k\}.$$

We can use a computer algebra package to solve equations (7.7)-(7.9) for a_0, b_0 and c_0 and then substitute these values into the equation $a_0^2 + b_0^2 = 1$ to obtain an equation of the form $P = 0$ where $P = \sum_{i=0}^8 t_i z_0^i$, $t_i \in \overline{\mathbb{Q}(p|_{G'}, q|_{G'}, x_0, y_0)}$ and $G' = G - v_0$. Since G' is infinitesimally VR-rigid by Theorem 7.1, we have $\overline{\mathbb{Q}(q|_{G'})} = \overline{\mathbb{Q}(p|_{G'})}$ by Lemma 7.2, and hence $t_i \in \overline{\mathbb{Q}(p|_{G'}, x_0, y_0)}$. Since z_0 is algebraically independent over $\overline{\mathbb{Q}(p|_{G'}, x_0, y_0)}$, we must have $t_i = 0$ for $0 \leq i \leq 8$. We may use equations (7.10) to simplify the expression for t_8 to give

$$t_8 = -4(c_j - z_j)^4 (c_j + z_j)^4 (a_j a_k - b_j b_k + 1)(a_j a_k + b_j b_k - 1) = 0.$$

Since $(G', q|_{G'})$ is quasi-generic and in standard position with respect to v_1 , we have $(a_j a_k - b_j b_k + 1)(a_j a_k + b_j b_k - 1) \neq 0$. Hence $c_j = \pm z_j$. By reflecting (G, q) in the xy -plane, we may assume that $c_j = z_j$. The fact that (G, p) and (G, q) are VR-equivalent now gives $c_i = z_i$ for all $v_i \in V$. This in turn implies that the frameworks (G, \tilde{p}) and (G, \tilde{q}) obtained by projecting (G, p) and (G, q) onto the xy -plane are equivalent. Since quasi-generic realisations of 2-connected graphs on the unit circle are globally rigid, (G, \tilde{p}) and (G, \tilde{q}) are congruent on the unit circle. Since $\tilde{p}(v_1) = (0, 1) = \tilde{q}(v_1)$, either $\tilde{p} = \tilde{q}$, or \tilde{p} can be obtained from \tilde{q} by reflection in the y -axis. This implies that either $p = q$, or p can be obtained from q by reflection in the yz -plane. Hence (G, p) and (G, q) are VR-congruent.

It remains to consider the case when all vertices of G other than v_1 have degree at least three. In this case it is easy to see that $G - e$ is a 2-connected graph with at least $|V| + 1$ edges for some $e \in E$. We may now use induction to deduce that $(G - e, p)$ is globally VR-rigid and hence that (G, q) is VR-congruent to (G, p) . \square

Lemma 5.2 is just a restatement of Theorem 7.4.

8. CLOSING REMARKS

1) We can also consider frameworks on families of concentric cylinders. Let $(\mathcal{Y}_1, \mathcal{Y}_2, \dots, \mathcal{Y}_n)$ be an ordered family of (not necessarily distinct) concentric cylinders where $\mathcal{Y}_i = \{(x, y, z) \in \mathbb{R}^3 : x^2 + y^2 = r_i\}$ and $r = (r_1, \dots, r_n)$ is a vector of positive real numbers. We say that the family is *generic* if the radii r_i are algebraically independent over \mathbb{Q} . A *framework* on $(\mathcal{Y}_1, \mathcal{Y}_2, \dots, \mathcal{Y}_n)$ is an ordered pair (G, p) consisting of a graph G and a realisation p such that $p(v_i) \in \mathcal{Y}_i$ for all $v_i \in V$.

In [19] it was proved that a graph is generically rigid on a family of concentric cylinders if and only if it is generically rigid on the cylinder. We may adapt the proof of Theorem 1.2 to obtain a similar result for generic global rigidity on families of cylinders, but our proof only works for *generic* families of cylinders. Given a framework (G, p) in \mathbb{R}^3 with $p(v_i) = (x_i, y_i, z_i)$, we define the family of concentric cylinders induced by p to be $\mathcal{Y}^p = (\mathcal{Y}_1, \mathcal{Y}_2, \dots, \mathcal{Y}_n)$ where $\mathcal{Y}_i = \{(x, y, z) \in \mathbb{R}^3 : x^2 + y^2 = x_i^2 + y_i^2\}$.

Theorem 8.1. *Let (G, p) be a generic framework in \mathbb{R}^3 . Then (G, p) is globally rigid on the family \mathcal{Y}^p of concentric cylinders induced by p if and only if G is either a complete graph on at most four vertices or G is 2-connected and redundantly rigid on the cylinder.*

Proof. Necessity follows from [13]. The proof of sufficiency is similar to the proof of Theorem 1.2. Theorem 4.6 and [14, Lemma 10] (which tells us that, if G has a realisation with a maximum rank equilibrium stress on some family of cylinders, then every generic realisation of G on a generic family of cylinders will have the same property) imply that (G, p) has a maximum rank equilibrium stress on \mathcal{Y}^p when G is an $\mathcal{M}_{2,2}^*$ -circuit. Combining this with [14, Theorem 9] (which tells us that (G, p) is globally rigid on \mathcal{Y}^p when it has a maximum rank equilibrium stress) implies that (G, p) is globally rigid on \mathcal{Y}^p when G is an $\mathcal{M}_{2,2}^*$ -circuit. The extension from a circuit to a 2-connected redundantly rigid graph can be accomplished using an ear decomposition of the rigidity matroid as in the proof of Theorem 1.2. \square

2) As noted in the proof of Theorem 8.1, [14, Theorem 9] tells us that a generic framework on a generic family of concentric cylinders is globally rigid if it has a maximum rank equilibrium stress. Here we show that the same result holds for the cylinder.

Theorem 8.2. *Let (G, p) be a generic framework on \mathcal{Y} with a maximum rank equilibrium stress (ω, λ) . Then (G, p) is globally rigid on \mathcal{Y} .*

Proof. The result will follow from Lemmas 5.1 and 5.2 if we can show that G is 2-connected and has at least $n + 1$ edges, where $n = |V(G)|$.

We first show that G is 2-connected. Suppose not, then $G = G_1 \cup G_2$ for subgraphs $G_i = (V_i, E_i)$ with $|V_1 \cap V_2| \leq 1$. Let $\omega^i = \omega|_{E_i}$. Then ω^i can be extended to a unique equilibrium stress (ω^i, λ^i) for $(G_i, p|_{G_i})$ on \mathcal{Y} , for $i = 1, 2$. We will consider the case when $|V_1 \cap V_2| = 1$ (the case when $V_1 \cap V_2 = \emptyset$ is simpler). Let $V_1 = \{v_1, v_2, \dots, v_k\}$ and $V_2 = \{v_k, v_{k+1}, \dots, v_n\}$. Then

$$\Omega(\omega) = \begin{pmatrix} A_1 & a^T & 0 \\ a & b & c \\ 0 & c^T & A_2 \end{pmatrix}, \text{ where } \Omega(\omega_1) = \begin{pmatrix} A_1 & a^T \\ a & b_1 \end{pmatrix}, \Omega(\omega_2) = \begin{pmatrix} b_2 & c \\ c^T & A_2 \end{pmatrix}$$

and $b_1 + b_2 = b$.

Since $\Omega(\omega_1)$ has column sums equal to zero, we can use elementary row operations to reduce $\Omega(\omega)$ to the form

$$\begin{pmatrix} A_1 & a^T & 0 \\ 0 & b_2 & c \\ 0 & c^T & A_2 \end{pmatrix}.$$

Since $\text{rank}(A_1 \ a^T) = \text{rank} \Omega(\omega_1) \leq |V_1| - 2$ and $\text{rank} \begin{pmatrix} b_2 & c \\ c^T & A_2 \end{pmatrix} = \text{rank} \Omega(\omega_2) \leq |V_2| - 2$, we have $\text{rank} \Omega(\omega) \leq n - 3$. This contradicts the hypothesis that (ω, λ) is a maximum rank equilibrium stress.

We next show G has at least $n + 1$ edges. In fact we will show the stronger fact that G is rigid on \mathcal{Y} . Since deleting an edge e with $\omega_e = 0$ will not reduce the rank of $\Omega_{\text{cyl}}(\omega, \lambda)$, we may assume that $\omega_e \neq 0$ for all edges e of G . Then every edge of G is contained in a $\mathcal{M}_{2,2}^*$ -circuit. Since G is connected, $\mathcal{M}_{2,2}^*$ -circuits are rigid, and the union of any two rigid subgraphs with a non-empty intersection is rigid, G is rigid on \mathcal{Y} . \square

3) We conjecture that the converse to Theorem 8.2 holds. That is, a generic framework (G, p) which is globally rigid on \mathcal{Y} must have a maximum rank equilibrium stress. This would be an analogue of the main result of [8]. It follows from our work that this conjecture holds when G is an $\mathcal{M}_{2,2}^*$ -circuit since we have shown that every $\mathcal{M}_{2,2}^*$ -circuit is both globally rigid and has a maximum rank equilibrium stress. One way to extend this to all graphs would be to prove a version of Lemma 5.6 showing that the ‘glueing’ operation preserves the property of having a maximum rank equilibrium stress. We can adapt the proof of Connelly [5, Lemma 10] to prove the special case when G_1 and G_2 have exactly two vertices in common but the general case seems challenging. An alternative approach to this problem would be to try to extend the recursive characterisation of $\mathcal{M}_{2,2}^*$ -circuits to a recursive characterisation of graphs which are 2-connected and redundantly rigid on \mathcal{Y} , and then use it to show that all generic realisations of these graphs have a maximum rank equilibrium stress.

4) In [21] frameworks on ‘expanding’ spheres were considered. Our v -free rigidity model is analogous to this for the cylinder. We believe that a similar technique to the proof of Theorem 6.3 could be used to extend our characterisation of v -free rigidity on \mathcal{Y} to the case when a subset of the vertices is allowed to move off the cylinder in a coordinated way. Since the analogous problem when all vertices are free to move off the cylinder independently is exactly the well known 3-dimensional rigidity problem, it may be of interest to extend this further to consider the case when two or more subsets of the vertex set can move independently off the cylinder. It is not hard to adapt the examples given in [21] to show that the obvious sparsity count is not sufficient to imply minimal rigidity when there are at least three subsets moving independently, but it may be sufficient when there are at most two subsets.

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APPENDIX: BASE GRAPHS FOR THE PROOFS OF THEOREMS 4.7 AND 6.1

We define a framework (G, p) for $G \in \{K_5^-, H_1, H_2\}$ which is infinitesimally rigid on \mathcal{Y} and has a nowhere zero maximum rank equilibrium stress (ω, λ) . These constructions form the base case for the inductive proofs of Theorems 4.7 and 6.1. We will use the labeling of the vertices given in Figure 2 and adopt the convention that ω_{ij} is the weight on the edge $v_i v_j$ in ω and λ_i is the weight on the vertex v_i in λ .

Case 1: $G = K_5^-$. Let (G, p) and (ω, λ) be defined by $p(v_1) = (0, 1, 0), p(v_2) = (1, 0, -1), p(v_3) = (\frac{1}{\sqrt{2}}, -\frac{1}{\sqrt{2}}, \frac{1}{3}), p(v_4) = (-1, 0, -\frac{1}{3}), p(v_5) = (\frac{1}{\sqrt{2}}, \frac{1}{\sqrt{2}}, \frac{1}{2})$,

$$(\omega_{12}, \omega_{13}, \omega_{14}, \omega_{15}, \omega_{23}, \omega_{24}, \omega_{25}, \omega_{34}, \omega_{45}) = (239, -216 - 654\sqrt{2}, 201 + 270\sqrt{2}, 756 + 616\sqrt{2}, 108 + 327\sqrt{2}, -\frac{1635}{2} - 852\sqrt{2}, 108 + 88\sqrt{2}, -108 - 327\sqrt{2}, -648 - 528\sqrt{2})$$

and

$$(\lambda_1, \lambda_2, \lambda_3, \lambda_4, \lambda_5) = (290 + 254\sqrt{2}, 1595 + 1397\sqrt{2}, 1524 + 870\sqrt{2}, 3045 + 2667\sqrt{2}, 1016 + 580\sqrt{2}).$$

It is straightforward to check that $\text{rank } R_{\text{cyl}}(G, p) = 13$, that $(\omega, \lambda) \cdot R_{\text{cyl}}(G, p) = 0$ and that $\text{rank } \Omega_{\text{cyl}}(\omega, \lambda) = 9$.

Case 2: $G = H_1$. Let (G, p) and (ω, λ) be defined by $p(v_1) = (0, 1, 0), p(v_2) = (-1, 0, -\frac{1}{3}), p(v_3) = (\frac{1}{\sqrt{2}}, -\frac{1}{\sqrt{2}}, \frac{1}{3}), p(v_4) = (1, 0, -1), p(v_5) = (0, -1, \frac{2}{3}), p(v_6) = (\frac{1}{\sqrt{2}}, \frac{1}{\sqrt{2}}, \frac{1}{2})$,

$$(\omega_{12}, \omega_{13}, \omega_{14}, \omega_{15}, \omega_{16}, \omega_{23}, \omega_{24}, \omega_{34}, \omega_{45}, \omega_{46}, \omega_{56}) = (1, 2\sqrt{2}, -\frac{361}{441} + \frac{10}{49}\sqrt{2}, -\frac{62}{147} - \frac{82}{147}\sqrt{2}, -\frac{20}{49} - \frac{80}{441}, \sqrt{2}, \frac{1}{2} + \sqrt{2}, -\sqrt{2}, \frac{32}{147} + \frac{12}{49}\sqrt{2}, \frac{4}{49} + \frac{16}{441}\sqrt{2}, -\frac{24}{49} - \frac{32}{147}\sqrt{2})$$

and

$$(\lambda_1, \lambda_2, \lambda_3, \lambda_4, \lambda_5, \lambda_6) = (-\frac{10}{9} - \frac{10}{9}\sqrt{2}, -3 - 3\sqrt{2}, -4 - 2\sqrt{2}, -\frac{13}{9} - \frac{13}{9}\sqrt{2}, \frac{4}{3} + \frac{4}{3}\sqrt{2}, \frac{8}{9} + \frac{4}{9}\sqrt{2}).$$

It is straightforward to check that $\text{rank } R_{\text{cyl}}(G, p) = 16$, that $(\omega, \lambda) \cdot R_{\text{cyl}}(G, p) = 0$ and that $\text{rank } \Omega_{\text{cyl}}(\omega, \lambda) = 12$.

Case 3: $G = H_2$. Let (G, p) and (ω, λ) be defined by $p(v_1) = (0, 1, 0), p(v_2) = (-1, 0, -\frac{1}{3}), p(v_3) = (\frac{1}{\sqrt{2}}, -\frac{1}{\sqrt{2}}, \frac{1}{3}), p(v_4) = (1, 0, -1), p(v_5) = (0, -1, \frac{2}{3}), p(v_6) = (\frac{1}{\sqrt{2}}, \frac{1}{\sqrt{2}}, \frac{1}{2}), p(v_7) = (-\frac{1}{\sqrt{2}}, -\frac{1}{\sqrt{2}}, -\frac{1}{4})$

$$(\omega_{12}, \omega_{13}, \omega_{14}, \omega_{17}, \omega_{23}, \omega_{24}, \omega_{34}, \omega_{45}, \omega_{46}, \omega_{47}, \omega_{56}, \omega_{57}, \omega_{67}) = (1, 2\sqrt{2}, \frac{4}{7}\sqrt{2} - \frac{13}{21}, \frac{8}{7} + \frac{8}{21}\sqrt{2}, \sqrt{2}, \frac{1}{2} + \sqrt{2}, -\sqrt{2}, -\frac{2652}{25165} - \frac{338}{25165}\sqrt{2}, -\frac{2764}{25165}\sqrt{2} - \frac{4652}{25165}, \frac{28424}{75495}\sqrt{2} + \frac{24784}{25165}, \frac{2764}{25165}\sqrt{2} + \frac{4652}{25165}, \frac{568}{3595} + \frac{16}{3595}\sqrt{2}, \frac{18608}{45297} + \frac{11056}{45297}\sqrt{2})$$

and

$$(\lambda_1, \lambda_2, \lambda_3, \lambda_4, \lambda_5, \lambda_6, \lambda_7) = (-\frac{74}{21}\sqrt{2} - \frac{82}{21}, -3 - 3\sqrt{2}, -2\sqrt{2} - 4, \frac{269}{105} - \frac{253}{105}\sqrt{2}, -\frac{4}{35}\sqrt{2} - \frac{12}{35}, -\frac{212}{315}\sqrt{2} - \frac{328}{315}, -\frac{704}{315}\sqrt{2} - \frac{1216}{315}).$$

It is straightforward to check that $\text{rank } R_{\text{cyl}}(G, p) = 19$, that $(\omega, \lambda) \cdot R_{\text{cyl}}(G, p) = 0$ and that $\text{rank } \Omega_{\text{cyl}}(\omega, \lambda) = 15$.

SCHOOL OF MATHEMATICAL SCIENCES, QUEEN MARY, UNIVERSITY OF LONDON, E1 4NS, U.K.

E-mail address: b.jackson@qmul.ac.uk

DEPARTMENT OF MATHEMATICS AND STATISTICS, LANCASTER UNIVERSITY, LA1 4YF, U.K.

E-mail address: a.nixon@lancaster.ac.uk