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Optimal Power Allocation for Multiuser Secure Communication in Cooperative Relaying Networks

Trung Q. Duong, Tiep M. Hoang, Chinmoy Kundu, Maged Elkashlan, and Arumugam Nallanathan

Abstract—We consider a cooperative relaying network in which a source communicates with a group of users in the presence of one eavesdropper. We assume that there are no source-user links and the group of users receive only retransmitted signal from the relay. Whereas, the eavesdropper receives both the original and retransmitted signals. Under these assumptions, we exploit the user selection technique to enhance the secure performance. We first find the optimal power allocation strategy when the source has the full channel state information (CSI) of all links. We then evaluate the security level through: i) ergodic secrecy rate and ii) secrecy outage probability when having only the statistical knowledge of CSIs.

Index Terms—Physical layer security, secrecy capacity, decode-and-forward, best user selection.

I. INTRODUCTION

Physical layer security has emerged and attracted the attention of communication society widely over the past decades. To deal with the leakage effect, many efforts to guarantee the security of cooperative networks have been made so far. Particularly, selection techniques are studied for the purpose of selecting suitable nodes in order to perform effective and safe transmission. Such selection techniques can be relay selection [1], jamming selection [2], [3], antenna selection [4], and user selection [5]. In general, such diverse selection techniques have been widely studied in literature.

Among the aforementioned techniques, user selection is an appropriate scheme for multiuser communications [5]. It exploits cooperative users to capture the cooperative diversity from a user group and enhance system security. We thus employ the user selection technique for choosing the best link among all links from a single relay to multiple users. Compared to closely previous related works, the differences are major, for example [5] exploited the user selection technique but did not discuss the source-eavesdropper link; while [6] did not discuss the impact of multiple users although the model is similar to ours. In contrast, we consider a severe scenario with the existence of source-eavesdropper link but no source-destination link. As such, the destination is in a much weaker position than the eavesdropper because the destination is receiving only one signal from the relay. We are thus motivated to examine such a secure cooperative network to understand the impact of the direct eavesdropping link on the secure performance. First, with the knowledge of instantaneous channel state information (CSI) of every channel, we find the optimal instantaneous source power which maximizes the instantaneous secrecy rate (SR). Then, only with the statistical knowledge of CSI, we evaluate the security level through the ergodic SR and secrecy outage probability (SOP).

II. SYSTEM MODEL AND USER SELECTION CRITERION

A. System Model

We consider the multiuser secure communication in cooperative relay networks, in which there are one source, one relay, one eavesdropper and multiple destinations, and all the nodes are single-antenna devices operating in the half-duplex mode. For notational simplicity, we denote the source, the relay, the k-th destination and the eavesdropper by S, R, kD (or simply kW) and E respectively. We assume that there is no direct link between S and kD due to loss and shadowing. In the first time slot, S broadcasts its signal while R attempts to decode the source signal. In the second time slot, R forwards the decoded signal to the best destination. We note that the signals transmitted by S and R are intercepted by E. Moreover, we assume that the CSI of kD as well as that of E is available at S and R. Hence, the availability of the CSI of E is a commonly used assumption when E is an active user of the system and not the intended destination for confidential messages [3], [4].

The channel between X ∈ {S, R} and Y ∈ {R, kW, E} (with Y ≠ X) is assumed to suffer from block Rayleigh fading with the channel gain $h_{XY} \sim CN(0, \Omega_{XY})$. Let $\gamma_{XY}$ denote the instantaneous signal-to-noise ratio (SNR) at Y for the signal transmitted by X. Then, we have $\gamma_{XY} = \frac{P_X}{N_Y} |h_{XY}|^2 \sim \text{Exp} \left( \frac{P_X}{N_Y} \Omega_{XY} \right)^2$ with $P_X$ being the transmit power at X and $N_Y$ being the additive white Gaussian noise at Y. By normalizing $N_Y$ (i.e., $N_Y = 0$ dB), $P_X$ can be understood as a replacement for $\frac{\Omega_{XY}}{N_Y}$. As such, in the first time slot, we have $\gamma_{SR} \sim \text{Exp} (\tau_{SR})$ and $\gamma_{SE} \sim \text{Exp} (\tau_{SE})$ where $\tau_{SR} = \frac{P_S \Omega_{SR}}{P_R}$ and $\tau_{SE} = \frac{P_S \Omega_{SE}}{P_R}$. Similarly, in the second time slot, we have $\gamma_{RK} \sim \text{Exp} (\tau_{RD})$ and $\gamma_{RE} \sim \text{Exp} (\tau_{RE})$ where $\tau_{RD} = \frac{P_R \Omega_{RD}}{P_R}$ and $\tau_{RE} = \frac{P_R \Omega_{RE}}{P_R}$.

B. User Selection Criterion

In this paper, the best user link selection is proposed, in which the destination corresponding to $\max_k \{ \gamma_{RK} \}$ is selected.

$CN(0, \Omega)$ denotes a complex Gaussian variable with zero-mean and variance $\Omega$. $\text{Exp} (m)$ denotes the exponential distribution with mean $m$. 
for receiving the signal retransmitted by R. We use $k^*\text{D}$ to denote the selected/strongest D corresponding to the index $k^* = \arg\max_k \gamma_{Rk}$. The instantaneous SNR for the R-$k^*$D channel is therefore $\gamma_{Rk^*} = \max_k \gamma_{Rk}$.

**Proposition 1.** The CDF and PDF of $\gamma_{Rk^*}$ can be, respectively, given by

$$F_{\gamma_{Rk^*}}(\gamma) = (1 - e^{-\gamma/T_RD})^K = 1 - \sum_k e^{-k(\gamma/T_RD)} \gamma,$$

$$f_{\gamma_{Rk^*}}(\gamma) = dF_{\gamma_{Rk^*}}(\gamma)/d\gamma = \sum_k (k/T_RD) e^{-k(\gamma/T_RD)} \gamma.$$

where $\sum_k (\gamma) \triangleq \sum_{k=1}^K (\gamma^k) (-1)^{k-1}$.

III. INSTANTANEOUS SOURCE POWER ALLOCATION STRATEGY

In this section, we examine the instantaneous source power which maximizes the SR. It should also be noted that the SR is the difference between the capacity of the desired link and that which maximizes the SR. It should also be noted that the SR is $R$, the end-to-end capacity of the channel from S to E with the help of R is given by

$$C_{k^*} = (1/2) \log_2 (1 + \max\{\gamma_{SR}, \gamma_{Rk^*}\}).$$

While E receives both versions of $x$, the end-to-end capacity of the channel from S to E with the help of R is given by

$$C_E = (1/2) \log_2 (1 + \max\{\gamma_{SR}, \gamma_{SE} + \gamma_{RE}\}).$$

For the sake of convenience, we let $X \triangleq \min\{\gamma_{SR}, \gamma_{Rk^*}\}$ and $Y \triangleq \min\{\gamma_{SR}, \gamma_{SE} + \gamma_{RE}\}$.

We assume that CSIs of all channels are perfectly known (full CSIs). The instantaneous SR of the system can be defined as [6]

$$C_D = |C_{k^*} - C_E| = (1/2) \log_2 ((1 + X)/(1 + Y)),$$

where $|x^+| = \max\{0, x\}$. We now consider finding the optimal transmit power $P_S$ that maximizes the SR given the maximum transmit power $P_{max}$ at S. Thus our SR maximization problem can be formulated as follows:

**P1** maximize $U(P_S) \triangleq (1 + X)/(1 + Y)$ subject to $0 < P_R \leq P_{max}$,

$X = \min\{\gamma_{SR}, \gamma_{Rk^*}\}$,

$Y = \min\{\gamma_{SR}, \gamma_{SE} + \gamma_{RE}\}$.

Let us label event $X$ as $X = \gamma_{SR} \Leftrightarrow |h_{SR}|^2 P_S \leq |h_{RE}|^2 P_R$, event $Y$ as $Y = \gamma_{SR} \Leftrightarrow |h_{RE}|^2 P_R \geq \max(|h_{SR}|^2 - |h_{SE}|^2, 0)$, and event $Z$ as $|h_{SR}|^2 \leq |h_{SE}|^2$. Respectively. In this case, $X$, $Y$, and $Z$ denote events $X = \gamma_{Rk^*}$, $Y = \gamma_{SE} + \gamma_{RE}$, and $|h_{SR}|^2 > |h_{SE}|^2$ respectively. Likewise, we label event $M$ as $P_{max} \geq |h_{SR}|^2 P_{max} \triangleq P_M$ and event $N$ as $P_{max} \geq |h_{SR}|^2 P_{max} \triangleq P_N$ respectively. While $M$, $N$ denote events $P_{max} < P_M$ and $P_{max} < P_N$ respectively.

Table I

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<tr>
<th>$M$</th>
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<td>$X'$</td>
<td>$0 \leq P_S \leq P_{max}$</td>
<td>$0 \leq P_S \leq P_{max}$</td>
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<tr>
<td>$X$</td>
<td>$0 \leq P_M &lt; P_{max}$</td>
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Table II

<table>
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<th>$Z$</th>
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<td>$Z$ and $N$</td>
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1) Scenario 1 ($X' \cap Y$): The objective function becomes $U(P_S) = 1 + \gamma_{SR} + |h_{SR}|^2$, which leads to $C_D = 0$ regardless of $P_S$ and $P_R$. It reveals that if any transmission is performed, it is not effective. Thus, S should not transmit signal, i.e., $P_S^{(1)} = 0$.

2) Scenario 2 ($X \cap Y$): The objective function becomes $U(P_S) = 1 + \gamma_{SR} + |h_{SR}|^2$. The first order derivative of $U(P_S)$ is then shown as $\frac{dU(P_S)}{dP_S} = \frac{\gamma_{SR} P_S}{1 + \gamma_{SR} + |h_{SR}|^2}$. Thus, $U(P_S)$ is a monotonically non-decreasing function when

$$|h_{SE}|^2 - |h_{SR}|^2 / |h_{SR}|^2 \leq |h_{SE}|^2 \leq |h_{SR}|^2.$$
Finally, the constraint $0 < P_S \leq P_{\text{max}}$ of (P1) is equivalent to $P_S \in \{ P_S^{(i)} \}_{i=1,2,3,4}$, i.e.,

$$\text{(P1)} \quad \text{maximize} \quad 0 < P_S \leq P_{\text{max}} \quad U(P_S)$$

$$\Leftrightarrow \text{(P2)} \quad \text{maximize} \quad P_S \in \{ P_S^{(i)} \}_{i=1,2,3,4} \quad U(P_S).$$

As such, (P2) can be readily solved by choosing the optimal $P_S^{(i)}, i \in \{1,2,3,4\}$, so that $U(P_S^{(i)})$ is maximal.

### IV. Performance Analysis

In this section, we evaluate the secure performance of our system model through two metrics: i) ergodic SR and ii) SOP. These two metrics are derived without requiring the knowledge of the instantaneous CSIs.

#### A. Ergodic Secrecy Rate

The ergodic SR in bits/s/Hz is given by

$$\langle C \rangle = (2 \ln 2)^{-1} \mathbb{E}_{\gamma_{SR}} \left\{ E_{Y|\gamma_{SR}} \left[ \omega_1(y, \gamma) \mid \gamma_{SR} = \gamma \right] \right\}$$

$$= (2 \ln 2)^{-1} \mathbb{E}_{\gamma_{SR}} \left\{ \omega_2(\gamma) \right\} \quad \text{(7)}$$

where

$$\omega_1(y, \gamma) \triangleq \int_0^\infty \ln \left( \frac{1+x}{1+y} \right) f_{X|\gamma_{SR}}(x|\gamma) dx,$$

$$\omega_2(\gamma) \triangleq \int_0^\infty \omega_1(y, \gamma) f_{Y|\gamma_{SR}}(y|\gamma) dy. \quad \text{(8)}$$

The second equality of (7) follows from that $X|\gamma_{SR}$ and $Y|\gamma_{SR}$ are independent. Their distributions are respectively shown in Propositions 2 and 3.

**Proposition 2.** The CDF of $X|\gamma_{SR} = \min \{ \gamma_{SR}, \gamma_{R_k} \} = \gamma$ is given by

$$F_{X|\gamma_{SR}}(x) = \begin{cases} F_{\gamma_{R_k}}(x), & \text{if } x < \gamma \\ 1, & \text{if } x \geq \gamma \end{cases}. \quad \text{(10)}$$

It is noted that $X|\gamma_{SR}$ is a mixed random variable, thus its PDF can be calculated as [10, Chapter 3]

$$f_{X|\gamma_{SR}}(x) = f_{\gamma_{R_k}}(x) + \left[ 1 - f_{\gamma_{R_k}}(\gamma) \right] \delta(x-\gamma), \quad \text{(11)}$$

if $x \leq \gamma$, where $\delta(x-\gamma)$ is a Dirac delta function in $x$.

**Proposition 3.** Let $Y_0 = \gamma_{SE} + \gamma_{RE}$, then the CDF and PDF of $Y_0$ can be, respectively, given by

$$F_{Y_0}(y) = 1 - \frac{\tau_{SE} - \tau_{RE}}{\tau_{SE} - \tau_{RE}} e^{-y/\tau_{SE} - \tau_{RE}}, \quad \text{(12)}$$

$$f_{Y_0}(y) = \left( e^{-y/\tau_{SE}} - e^{-y/\tau_{RE}} \right) \left( \tau_{SE} - \tau_{RE} \right)^{-1}. \quad \text{(13)}$$

Let $Y|\gamma_{SR} = \min \{ \gamma_{SR}, Y_0 \} |\gamma_{SR} = \gamma$, we can obtain the CDF and the PDF of $Y|\gamma_{SR}$ by replacing $\gamma_{R_k}$ in (10)-(11) with $Y_0$.

Substituting (11) into (8) and using the sifting property of Dirac delta function, we have

$$\omega_1(y, \gamma) = \sum_k e^{k \gamma_{RD}} \left[ E_1 \left( \frac{k(1+y)}{\gamma_{RD}} \right) - E_1 \left( \frac{k(1+\gamma)}{\gamma_{RD}} \right) \right]. \quad \text{(14)}$$

where $E_1(x) = \int_{y}^{\infty} \frac{e^{-u}}{u} du$ for $y \leq \gamma$; otherwise, $\omega_1(y, \gamma) = 0$ for $y > \gamma$. As such, $\omega_2(\gamma)$ in (9) can be reduced to $\omega_2(\gamma) = \int_0^\gamma \omega_1(y, \gamma) f_{Y_0}(y) dy$ by integrating over $y \leq \gamma$. Using (13) and (14) to evaluate $\omega_2(\gamma)$ again, we then arrive at

$$\omega_2(\gamma) = \sum_k e^{k \gamma_{RD}} \left( \tau_{SE} - \tau_{RE} \right)^{-1}$$

$$\times \left\{ \int_0^\gamma E_1 \left( k(1+y)/\tau_{RD} \right) \left( e^{-y/\tau_{SE}} - e^{-y/\tau_{RE}} \right) dy + E_1 \left( k(1+\gamma)/\tau_{RD} \right) \left( \tau_{SE}e^{-y/\tau_{SE}} - \tau_{RE}e^{-y/\tau_{RE}} \right) \right. \right.

$$\left. - E_1 \left( k(1+\gamma)/\tau_{RD} \right) \left( \tau_{SE} - \tau_{RE} \right) \right\}. \quad \text{(15)}$$

Finally, substituting (15) into (7) and using [11, Eq. (5.231.2)], we obtain the ergodic SR as shown at the bottom of the next page.

#### B. Secrecy Outage Probability

The SOP is given by

$$P_{\text{out}}(\zeta) = P \left\{ C_{\Delta} \leq R \right\} = \int_0^\infty P(\gamma, \zeta)f_{Y|\gamma_{SR}}(\gamma)d\gamma \quad \text{(17)}$$

where $\zeta = 2^{2R} - 1$ and the term $P(\gamma, \zeta)$ is defined as

$$P(\gamma, \zeta) \triangleq P \left\{ X - (\zeta+1)Y \leq \zeta |\gamma_{SR} = \gamma \right\}$$

$$= \int_0^\infty F_{X|\gamma_{SR}}(\gamma(1+y) + \zeta) \times \left[ f_{Y_0}(y) + (1 - f_{Y_0}(\gamma)) \right] dy \quad \text{(18)}$$

where $f_{Y|\gamma_{SR}}(y|\gamma)$, along with $f_{Y_0}(y)$, is presented in Proposition 3. Combining (18) with (10), we then have two cases:

- If $0 \leq \gamma \leq \zeta$, we deduce from (10) that $F_{X|\gamma_{SR}}((\zeta+1)y+\zeta) = 1$ and hence (18) becomes

$$P(\gamma, \zeta) = f_{Y_0}(\gamma) + (1 - f_{Y_0}(\gamma)). \quad \text{(19)}$$

- If $\zeta < \gamma < \infty$, we deduce from (10) that $F_{X|\gamma_{SR}}((\zeta+1)y+\zeta) = f_{Y_0}(\gamma) + (1 - f_{Y_0}(\gamma))$ if $y < \frac{\zeta}{\gamma}$, and $F_{X|\gamma_{SR}}((\zeta+1)y+\zeta) = 1$ if $y \geq \frac{\zeta}{\gamma}$. Hence (18) becomes

$$P(\gamma, \zeta) = J(\gamma, \zeta) + 1 - F_{Y_0} \left( \frac{\zeta-\gamma}{\zeta+1} \right) \quad \text{(20)}$$

**Note:** The above equations are derived for the case when $0 < P_S \leq P_{\text{max}}$. For other values of $P_S$, different metrics and calculations may be required. The above equations are valid under the assumption that $P_S$ is bounded by $P_{\text{max}}$.
where

\[
J(\gamma, \zeta) = \left[ \gamma_{SE} - \gamma_{RE} \right]^{-1} \sum_k e^{-\frac{k}{\gamma_{RD}}} \times \left\{ 1 - e^{-\left( \frac{k(\gamma - \zeta)}{\gamma_{RD}} + \frac{k(\gamma - \gamma_{SR})}{(\gamma_{RD} + 1)} \right)} \left[ k(\zeta + 1) \gamma_{RD} + 1 \right]^{-1} \right\}^{-1} \left[ k(\zeta + 1) \gamma_{RD} + 1 \right]^{-1} \right\}^{-1}
\]

with \( \sum_k \gamma(\cdot) \triangleq \sum_{k=0}^{K} (K)(-1)^k(\cdot) \).

Proceeding to analyze the integral (17) with the help of (19)–(20), we can rewrite (17) as

\[
P_{out}(\zeta) = 1 - I_1(\zeta) + I_2(\zeta)
\]

where

\[
I_1(\zeta) = \gamma_{SR}^{2\gamma_{SE}} \frac{e^{-\frac{\gamma_{SR}}{\gamma_{SE}}}}{\gamma_{SR}(\gamma_{SE} - \gamma_{RE})} \left[ \gamma_{SR} + (\zeta + 1) \gamma_{SE} + (\zeta + 1) \gamma_{RE} \right],
\]

\[
I_2(\zeta) = \gamma_{SR}^{2\gamma_{SE}} \frac{e^{-\frac{\gamma_{SR}}{\gamma_{SE}}}}{\gamma_{SR}(\gamma_{SE} - \gamma_{RE})} \left[ \gamma_{SR} - (k/\gamma_{RD} + 1/\gamma_{SE} + 1/\gamma_{SR}) \left[ k(\zeta + 1) / \gamma_{RD} + 1 / \gamma_{SE} \right]^{-1} \right] \times \left[ k(\zeta + 1) / \gamma_{RD} + 1 / \gamma_{SE} \right]^{-1} \left[ k(\zeta + 1) / \gamma_{RD} + 1 / \gamma_{RE} \right]^{-1} \times \left[ k(\zeta + 1) / \gamma_{RD} + 1 / \gamma_{RE} \right]^{-1} \times \left[ k(\zeta + 1) / \gamma_{RD} + 1 / \gamma_{RE} \right]^{-1}
\]

V. RESULTS

In this section, we provide some numerical results to validate analytical results. Unless specifically stated, we set parameters \( \frac{P_{max}}{N_0} = 40 \text{ dB} \) and \( \frac{P_0}{N_0} \approx 26.99 \text{ dB} \). Recall that \( \frac{P_0}{N_0} \) is considered as \( P_S \) due to \( N_0 \) being set to 0 dB.

\[
(\gamma) = (2\ln 2)^{-1} \left( \gamma_{SE} - \gamma_{RE} \right)^{-1} \sum_k e^{-\frac{k}{\gamma_{RD}}} \left\{ \gamma_{RE} \left[ e^{\gamma_{RE}} E_1(k/\gamma_{RD} + 1/\gamma_{SE} + 1/\gamma_{RE}) - E_1(k/\gamma_{RD} + 1/\gamma_{SR}) \right] \right. \left. - \gamma_{SE} \left[ e^{\gamma_{SE}} E_1(k/\gamma_{RD} + 1/\gamma_{SR} + 1/\gamma_{SE}) - E_1(k/\gamma_{RD} + 1/\gamma_{SR}) \right] \right\}.
\]

VI. CONCLUSIONS

In this work, we have considered a secured cooperative relaying network with the best user selection technique. We have proposed the optimal transmit power allocation strategy to help the instantaneous SR attain its maximum. Besides, we have investigated the secure performance of the system through the average SR and the SOP.

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